

Equazioni matriciali lineari con più termini nella risoluzione numerica di SPDEs

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$$A\mathbf{X} + \mathbf{X}B + D = 0$$

Eigenvalue pbs and tracking, Control, MOR, Assignment pbs, Riccati eqn

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Stability analysis in Control and Dynamical systems, Signal processing, eigenvalue computations

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Focus: All or some of the matrices are large (and possibly sparse)

Solving the Lyapunov equation. The problem Approximate ${\bf X}$ in:

 $A\mathbf{X} + \mathbf{X}A^{\top} + BB^{\top} = 0$ $A \in \mathbb{R}^{n \times n} \text{ neg.real} \qquad B \in \mathbb{R}^{n \times p}, \qquad 1 \le p \ll n$

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Time-invariant linear system:

$$\mathbf{x}'(t) = A\mathbf{x}(t) + B\mathbf{u}(t), \qquad \mathbf{x}(0) = x_0$$

Closed form solution:

$$\mathbf{X} = \int_0^\infty e^{-tA} B B^\top e^{-tA^\top} dt$$

 \Rightarrow X symmetric semidef.

see, e.g., Antoulas '05, Benner '06

Linear systems vs linear matrix equations

Large linear systems:

$$Ax = b, \quad A \in \mathbb{R}^{n \times n}$$

- Krylov subspace methods (CG, MINRES, GMRES, BiCGSTAB, etc.)
- \bullet Preconditioners: find P such that

$$AP^{-1}\widetilde{x} = b \qquad x = P^{-1}\widetilde{x}$$

is easier and fast to solve

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Large linear matrix equations:

$$A\mathbf{X} + \mathbf{X}A^{\top} + BB^{\top} = 0$$

- No preconditioning to preserve symmetry
- ${\bf X}$ is a large, dense matrix \Rightarrow low rank approximation

$$\mathbf{X} \approx \widetilde{X} = Z Z^{\top}, \quad Z \text{ tall}$$

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Large linear matrix equations:

$$A\mathbf{X} + \mathbf{X}A^{\top} + BB^{\top} = 0$$

Kronecker formulation:

$$(A \otimes I + I \otimes A)x = b$$
 $x = \operatorname{vec}(\mathbf{X})$

Projection-type methods

Given an approximation space \mathcal{K} ,

 $\mathbf{X} \approx X_m \quad \operatorname{col}(X_m) \in \mathcal{K}$ Galerkin condition: $R := AX_m + X_m A^\top + BB^\top \perp \mathcal{K}$ $V_m^\top R V_m = 0 \qquad \mathcal{K} = \operatorname{Range}(V_m)$

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Assume $V_m^{\top}V_m = I_m$ and let $X_m := V_m Y_m V_m^{\top}$. Projected Lyapunov equation:

$$V_m^{\top} (A V_m Y_m V_m^{\top} + V_m Y_m V_m^{\top} A^{\top} + B B^{\top}) V_m = 0$$

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$$V_m^{\top} (AV_m Y_m V_m^{\top} + V_m Y_m V_m^{\top} A^{\top} + BB^{\top}) V_m = 0$$

$$(V_m^{\top} A V_m) Y_m + Y_m (V_m^{\top} A^{\top} V_m) + V_m^{\top} BB^{\top} V_m = 0$$

Early contributions: Saad '90, Jaimoukha & Kasenally '94, for $\mathcal{K} = \mathcal{K}_m(A, B) = \text{Range}([B, AB, \dots, A^{m-1}B])$ More recent options as approximation space

Enrich space to decrease space dimension

• Extended Krylov subspace

$$\mathcal{K} = \mathcal{K}_m(A, B) + \mathcal{K}_m(A^{-1}, A^{-1}B),$$

that is, $\mathcal{K} = \text{Range}([B, A^{-1}B, AB, A^{-2}B, A^2B, A^{-3}B, \dots,])$ (Druskin & Knizhnerman '98, Simoncini '07) More recent options as approximation space

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• Rational Krylov subspace

$$\mathcal{K} = \mathbb{K} := \operatorname{Range}([B, (A - s_1 I)^{-1} B, \dots, (A - s_m I)^{-1} B])$$

usually, $\{s_1, \ldots, s_m\} \subset \mathbb{C}^+$ chosen a-priori

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In both cases, for $Range(V_m) = \mathcal{K}$, projected Lyapunov equation:

$$(V_m^{\top}AV_m)Y_m + Y_m(V_m^{\top}A^{\top}V_m) + V_m^{\top}BB^{\top}V_m = 0$$

 $X_m = V_m Y_m V_m^\top$

Multiterm linear matrix equation

$$A_1 \mathbf{X} B_1 + A_2 \mathbf{X} B_2 + \ldots + A_\ell \mathbf{X} B_\ell = C$$

Applications:

- Matrix least squares
- Control
- (Stochastic) PDEs
- ...

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Main device: Kronecker formulation

$$(B_1^{\top} \otimes A_1 + \ldots + B_{\ell}^{\top} \otimes A_{\ell}) x = c$$

Iterative methods: matrix-matrix multiplications and rank truncation (Benner, Breiten, Bouhamidi, Chehab, Damm, Grasedyck, Jbilou, Kressner, Matthies, Onwunta, Raydan, Stoll, Tobler, Zander, and many others...)

Matrix equations in PDEs

The Poisson equation - revisited

$$-u_{xx} - u_{yy} = f$$
, in $\Omega = (0, 1)^2$

+ Dirichlet b.c. (zero b.c. for simplicity)

Usual discretization \Rightarrow Au = b (with $A = T \otimes I + I \otimes T$)

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Usual discretization \Rightarrow Au = b (with $A = T \otimes I + I \otimes T$)

Discretization: $U_{i,j} \approx u_{x_i,y_j}$, with (x_i, y_j) interior nodes, so that h: meshsize

$$u_{xx}(x_i, y_j) \approx \frac{U_{i-1,j} - 2U_{i,j} + U_{i+1,j}}{h^2} = \frac{1}{h^2} [1, -2, 1] \begin{bmatrix} U_{i-1,j} \\ U_{i,j} \\ U_{i+1,j} \end{bmatrix}$$
$$u_{yy}(x_i, y_j) \approx \frac{U_{i,j-1} - 2U_{i,j} + U_{i,j+1}}{h^2} = \frac{1}{h^2} [U_{i,j-1}, U_{i,j}, U_{i,j+1}] \begin{bmatrix} 1 \\ -2 \\ 1 \end{bmatrix}$$

$$T\mathbf{U} + \mathbf{U}T = F, \qquad b = \operatorname{vec}(F)$$

 $-\Delta u = 1, \quad \Omega = (0,1)^3 \quad \Rightarrow \quad A = (T \otimes I \otimes I + I \otimes T \otimes I + I \otimes I \otimes T)$

 $-\Delta u = 1, \quad \Omega = (0,1)^3 \quad \Rightarrow \quad A = (T \otimes I \otimes I + I \otimes T \otimes I + I \otimes I \otimes T)$ CG for Ax = b vs lterative solver for $(I \otimes T + T \otimes I)\mathbf{U} + \mathbf{U}T = F$ $T \in \mathbb{R}^{n \times n}, A \in \mathbb{R}^{n^3 \times n^3}, \qquad n = 50$



A 3D convection-diffusion equation

 $-\epsilon\Delta u+\mathbf{w}\cdot\nabla u=1,$ in $\Omega=(0,1)^3,$ with convection term $\mathbf{w}=(x\sin x,y\cos y,e^{z^2-1})$

Sylvester equation:

$$[I \otimes (T_1 + \Phi_1 B_1) + (T_2 + \Psi_2 B_2)^\top \otimes I] \mathbf{U} + \mathbf{U} (T_3 + B_3 \Upsilon_3) = \mathbf{1} \mathbf{1}^\top$$

ϵ	n_x	FGMRES+AGMG	GMRES+MI20	Sylv Solver
		CPU time (# its)	CPU time (# its)	CPU time (# its)
0.0050	100	8.0207 (15)	9.7207 (7)	0.5677 (22)
0.0010	100	7.6815 (14)	9.4935 (7)	0.5446 (22)
0.0005	100	7.3914 (14)	9.6274 (7)	0.5927 (24)

- Also for more general, separable coeff., operators on uniform grids
- If not separable coeff., use as preconditioner

(Palitta & Simoncini (tr 2014))

... A classical approach

Matrix formulation is not new...

- Bickley & McNamee, 1960: Early literature on difference equations
- Wachspress, 1963: Model problem for ADI algorithm
- Ellner & Wachspress (1980's): interplay between the matrix and vector formulations (via preconditioning)

Novel solvers for matrix equations allow faster convergence

PDEs with random inputs

Stochastic steady-state diffusion eqn: Find $u: D \times \Omega \to \mathbb{R} \ s.t. \ \mathbb{P}$ -a.s.,

$$\begin{cases} -\nabla \cdot (a(\mathbf{x}, \omega) \nabla u(\mathbf{x}, \omega)) = f(\mathbf{x}) & in D \\ u(\mathbf{x}, \omega) = 0 & on \partial D \end{cases}$$

f: deterministic;

a: f. of finite no. of real-valued random variables $\xi_r : \Omega \to \Gamma_r \subset \mathbb{R}$ Common choice: truncated Karhunen–Loève (KL) expansion,

$$a(\mathbf{x},\omega) = \mu(\mathbf{x}) + \sigma \sum_{r=1}^{m} \sqrt{\lambda_r} \phi_r(\mathbf{x}) \xi_r(\omega),$$

 $\mu(\mathbf{x})$: expected value of diffusion coef. σ : std dev. $(\lambda_r, \phi_r(\mathbf{x}))$ eigs of the integral operator \mathcal{V} wrto $V(\mathbf{x}, \mathbf{x}') = \frac{1}{\sigma^2} C(\mathbf{x}, \mathbf{x}')$ $(\lambda_r \searrow C: D \times D \to \mathbb{R}$ covariance fun.)

Discretization by stochastic Galerkin

Approx with space in tensor product form $\mathcal{X}_h \times S_p$

$$\mathcal{A}\mathbf{x} = \mathbf{b}, \qquad \mathcal{A} = G_0 \otimes K_0 + \sum_{r=1}^m G_r \otimes K_r, \quad \mathbf{b} = \mathbf{g}_0 \otimes \mathbf{f}_0,$$

x: expansion coef. of approx to u in the tensor product basis $\{\varphi_i \psi_k\}$ $K_r \in \mathbb{R}^{n_x \times n_x}$, FE matrices (sym) $G_r \in \mathbb{R}^{n_\xi \times n_\xi}$, $r = 0, 1, \ldots, m$ Galerkin matrices associated w/ S_p (sym.)

 \mathbf{g}_0 : first column of G_0

 $\mathbf{f}_0:~\mathsf{FE}$ rhs of deterministic PDE

$$n_{\xi} = \dim(S_p) = \frac{(m+p)!}{m!p!}$$

^a S_p set of multivariate polyn of total degree $\leq p$

The matrix equation formulation

 $(G_0\otimes K_0+G_1\otimes K_1+\ldots+G_m\otimes K_m)\,\mathbf{x}=\mathbf{g}_0\otimes\mathbf{f}_0$ transforms into

 $K_0 \mathbf{X} G_0 + K_1 \mathbf{X} G_1 + \ldots + K_m \mathbf{X} G_m = F, \qquad F = \mathbf{f}_0 \mathbf{g}_0^\top$ $(G_0 = I)$

Solution strategy. Conjecture:

• $\{K_r\}$ from trunc'd Karhunen–Loève (KL) expansion

 $\label{eq:Seek} \begin{matrix} \Downarrow \\ \mathsf{Seek low rank } \widetilde{X} \approx \mathbf{X} \end{matrix}$

(Possibly extending results of Gradesyk, 2004)

Matrix Galerkin approximation of the deterministic part. 1

Approximation space \mathcal{K}_k and basis matrix V_k : $\mathbf{X} \approx X_k = V_k Y$

$$V_k^{\top} R_k = 0, \qquad R_k := K_0 X_k + K_1 X_k G_1 + \ldots + K_m X_k G_m - \mathbf{f}_0 \mathbf{g}_0^{\top}$$

Computational challenges:

- Generation of \mathcal{K}_k involved m+1 different matrices $\{K_r\}$!
- Matrices K_r have different spectral properties
- n_x, n_ξ so large that X_k, R_k should not be formed !

Ongoing joint project with Catherine Powell, David Silvester, Univ. Manchester

Matrix Galerkin approximation of the deterministic part. 2

Approximation space \mathcal{K}_k and basis matrix V_k : $\mathbf{X} \approx X_k = V_k Y$

 $V_k^{\top} R_k = 0, \qquad R_k := K_0 X_k + K_1 X_k G_1 + \ldots + K_m X_k G_m - \mathbf{f}_0 \mathbf{g}_0^{\top}$

• "Divide" by K_0 (stiffness matrix)

$$\widehat{\mathbf{X}} + \widehat{K}_1 \widehat{\mathbf{X}} G_1 + \ldots + \widehat{K}_m \widehat{\mathbf{X}} G_m = \widehat{\mathbf{f}}_0 \mathbf{g}_0^\top$$

• Shift matrices by α_r , $r = 1, \ldots, m$ to get similar spectral interval

$$\widehat{\mathbf{X}}\left(I - \sum_{r=1}^{m} \alpha_r G_r\right) + (\widehat{K}_1 + \alpha_1 I)\widehat{\mathbf{X}}G_1 + \ldots + (\widehat{K}_r + \alpha_m I)\widehat{\mathbf{X}}G_m = \widehat{\mathbf{f}}_0 \mathbf{g}_0^\top$$

- Generate $\mathcal{K}_k = \bigcup_r \mathbb{K}_k(\widehat{K}_r + \alpha_r I, \widehat{\mathbf{f}}_0)$
- Solve projected (reduced) problem by matrix-oriented CG

Example 1. SIFISS 1.0, Q1 Finite Elements

 $-\nabla \cdot (a\nabla u) = 1$, $D = (0,1)^2$, $u|_{\partial D} = 0$, (fast decay)^a $n_x = 65,025$ tol=10⁻⁵ (relative soln change)

^aEigel Gittelson Schwab Zander, CMAME 2014

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	m+1	p	n_{ξ}	$dim(\mathcal{K}_k)$	avr # inner	Elapsed Time
	9	3	165	139	14.60	184.19
	9	4	495	171	16.83	220.60
	9	5	1287	171	17.50	228.67
	11	3	286	168	14.60	232.48
	11	4	1001	168	15.80	232.94
	11	5	3003	213	18.00	295.10
	16	3	816	165	13.25	281.85
CG	9	5	1287		20	1229.3

(Direct solves for generating each rational space \mathbb{K}_k)

Dell PowerEdge R620 w/ 32 procs Intel(R) Xeon(R) CPU E5-2640 v2 @ 2.00GHz

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	9	4	495	167	16.83	62.19
	9	5	1287	167	17.50	63.73
	11	3	286	166	14.60	58.84
	11	4	1001	166	15.40	65.49
	11	5	3003	210	17.83	83.00
	16	3	816	166	13.25	80.31
CG	9	5	1287		20	1229.3

(Iterative solves for generating each rational space \mathbb{K}_k)

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Conclusions

Multiterm (Kron) linear equations is the new challenge

- Great advances in solving really large linear matrix equations
- Second order (matrix) challenges rely on strength and maturity of linear system solvers
- Low-rank tensor formats is the new generation of approximations

Reference for linear matrix equations:

* V. S., Computational methods for linear matrix equations,
March 2013, Submitted

available at www.dm.unibo.it/~simoncin