

# **Polygonal Mesh**

- Representation and properties
- Data Structure
- Simplification, compression, LOD
- Parametrization, Fairing



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# **Polygonal Mesh**



78K vertices 160K triangle



Polygonal meshes are simply collections of polygons, or "faces," that together form the "skin" of an object. They have become a standard way of representing a broad class of shapes in graphics.



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#### **Complex meshes in CG..**





### **Complex meshes in art...**



[Digital Michelangelo Project, 2000]



#### 2,000,000,000 faces

#### **Objectives:**

- rendering
- storage
- transmission
- scalability



## **Complex meshes in CAD**

- Submarine Auxiliary Machine Room
- 500,000 polygons



#### Courtesy General Dynamice, Electric Boat Div.

# Meshes in medical imaging...



#### Visible-body: 512x512x1734 punti





### **Complex meshes for terrain**



#### topography

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### erritorial Data: special meshes





#### 16K x 16K vertices ~537 milion of triangles

View-dependent, Geometric model + bump (texture) mapping

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### **Height field**

#### Visualize the explicit function: z = f(x,y)

# A height field is defined on a regular grid of quotes

$$h: [0, N-1]^2 \rightarrow \mathbf{R},$$

where N is huge

Store the height as an image (i.e. Format gif)





### **Geometry images**



#### Surface

$$(u, v, z(u, v))$$
$$(u_{ij}, v_{ij}) = (i\Delta u, j\Delta v)$$
$$Z_{ij} = z(u_{ij}, v_{ij})$$

#### **Geometry image**



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### .obj from google earth



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### Height field as a TIN

#### Alternatively, a height field is stored as triangular irregular networks (TIN) or mesh: a vector-based representation of a surface

use an optimal number of polygons to approximate a surface to a given level of detail and accuracy





## **Mesh generation**

Mesh from points:

triangulation or tetraedralization

- Sampling by a digitizer
- Reconstruction from multiple views
- 3D Scanner
- Territorial models
- Mesh from surfaces: tessellation (for surface rendering)
- Mesh from data volumes: polygonalization
  - Volumetric data (isosurfacing: extract a surface of constant value through the volume)



## **Mesh representation**

- Structured (regular) mesh : all internal vertices are surrounded by a costant number of elements.
- Semi-regular mesh is obtained by regular subdivision of an irregular mesh: all the vertices are regular except for a small number of extraordinary vertices





М

#### Mesh & Manifold

 $\mathcal{M}$  is a 2D-manifold of arbitrary topology embedded in  $\mathbb{R}^3$ 

**M** is a piecewise linear approximation of *M* 

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Μ



#### Manifold

A topological space in which every point has a neighborhood homeomorphic to  $\mathbb{R}^n$  (topological disc) is called an *n*-dimensional (or *n*-) manifold





2-manifold

Not a manifold

Earth is an example of a 2-manifold

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#### **Triangle Mesh**

A triangle mesh M consists of a **geometric** and a **topological** component, where the latter can be represented by a graph structure of the form (V,E,F) consisting of

Vertices

$$V = \{1, \dots, N_{v}\}$$

Edges

$$E = \left\{ (i, j) \in V \times V : X_j \in N(X_i) \right\}$$

between two vertices

Faces

$$F = \left\{ (i, j, k) \in V \times V \times V : (i, j), (i, k), (k, j) \in E \right\}$$
  
between edges



# **Approximation Quality**

 If a sufficiently smooth surface is approximated by a piecewise linear function, the approximation error is of the order O(h<sup>2</sup>), with h denoting the maximum edge length.



.....the error is reduced by a factor of about 1/4 ....

[Botsch et al.]

• The actual magnitude of the approximation error depends on the curvature of the underlying smooth surface.



- Boundaries of tangible physical objects are two-dimensional manifolds.
- They reside in (are embedded into, are subspaces of) the ambient three-dimensional Euclidean space.
- Two-dimensional manifolds are also called embedded surfaces (or simply surfaces).
- Can often be described by the map  $S: D \subset R^2 \to R^3$   $U \to u$   $D \subset R^2$  is a parameter domain.
  - the map  $S(u,v) = (x(u,v) \quad y(u,v) \quad z(u,v))^T$

is a global parameterization (embedding) of the manifold.

Smooth global parameterization does not always exist or is easy to find.

X



The description of most manifolds requires more than one parametric domain (more than one patch).
 A single patch is adequate for only the simplest manifolds.





### **Charts and atlases**

U

Chart  $\alpha : U_{\alpha} \to \mathbb{R}^2$ 

•X

 $\alpha(U_{\alpha}) \subset R^2$ 



A homeomorphism  $\alpha: U_{\alpha} \to R^2$ from a neighborhood  $U_{\alpha} \subset \mathcal{M}$ to  $R^2$  is called a **chart** 

It is not generally possible to describe a manifold with just one chart A collection of charts  $(U_{\alpha}, \alpha)_{\alpha \in A}$ whose domains cover the manifold is called an **atlas**  $\bigcup U_{\alpha} = \mathcal{M}$ 



#### **Charts and atlases**





## Manifolds with boundary



A topological space in which every point has an open neighborhood homeomorphic to either

• topological disc  $\mathbb{R}^n$ ; or

• topological half-disc[0,  $\infty$ ) ×  $\mathbb{R}^{n-1}$ 

is called a manifold with boundary

Points with disc-like neighborhood are called interior, denoted by Int(X) Points with half-disc-like neighborhood are called **boundary**, denoted by  $\partial X$ 



### Manifold meshes

A mesh is a 2-manifold if

 Neighborhood of each interior vertex is homeomorphic to a disc
 Neighborhood of each boundary vertex is homeomorphic to a half-disc







Two kinds of vertices allowed:

Internal

**Boundary** 

<sup>-</sup> Universita di Bologna



## Non-manifold meshes



A triangle mesh is a 2-manifold if it contains neither **non-manifold edges** nor **non-manifold vertices** nor **self-intersections** 

- non-manifold edge has more than two incident triangles
- non-manifold vertex is incident to more than one fan of triangles







- Discrete surface representation
- The mesh is a piece-wise planar approximation obtained by gluing the polygonal faces together,

#### Connectivity (Topological) data

The mesh is a purely **topological** object and does not contain any geometric properties **Geometric data** 

The **geometric realization** of the mesh is defined by specifying the coordinates of the vertices  $X_i \in R^3$  for all  $i \in V$ 

Orientation data



#### **Representing meshes**

- Geometry:
  - Vertex Coords
    - $(x_1, y_1, z_1)$  $(x_2, y_2, z_2)$
    - $(x_n, y_n, z_n)$
- Orientation
  - List of normals
- Connectivity
  - List of triangles
    - $(i_1, j_1, k_1)$  $(i_2, j_2, k_2)$
    - $(i_m, j_m, k_m)$



STI	normal	n <sub>x</sub>	n <sub>y</sub>	n <sub>z</sub>		vertex	х	У	Z
Fyamola	0	-1	0	-0-		0	0	0	0
	1	-0.707	0.707	0		1	1	0	0
	2	0.707	0.707	0	—	2	1	1	0
	3	1	0	0		3	0.5	1.5	0
	4	0	-1	0		4	0	1	0
	-	0	0	1		5	0	0	1
	5	0	0	-1		6	1		1
	6					-	1		1
$\mathbf{n}_1$ $\mathbf{n}_2$ $\mathbf{n}_2$						7	1		
						8	0.5	1.5	1
$_{8}$ $_{4}$ $\times$				face		9	0	1	1
2	, n <sub>6</sub>		fa			vertices	associated norma		rmal
			0	(left)		0,5,9,4	0,0,0,0		
	<b>n</b> <sub>2</sub>		1	(roof left)		3,4,9,8	1,1,1,	1	
$ \mathbf{n}_5 \bigtriangleup 5 $	~		2	(roof right)		2,3,8,7	2,2,2,2	2	
$\begin{array}{c c} & & & & \\ & & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & & \\ & & & \\$						1,2,7,6	3,3,3,3		
z 6			4	(bottom)		0,1,6,5	4,4,4,4	4	
			5	(front)		5,6,7,8,9	5,5,5,5	5,5	
4		A MATER ST	UDIC 6	(back)		0,4,3.2.1	6,6,6.6	6,6	



### **Normal Vectors**

Local connectivity
Valence of a vertex: number of incident edges at a vertex
Vertex 1-ring  $N(i) = \{j \in V : (i, j) \in E\} \subset V$ Face 1-ring  $N(f) = \{(i, j, k) \in F : j, k \in V\} \subset T$ Normal Computation

**Face Normal** 



$$f = (i, j, k) \in T, \quad n_f = \frac{(X_j - X_i) \times (X_k - X_i)}{\left\| (X_j - X_i) \times (X_k - X_i) \right\|}$$
  
Vertex Normal

$$\forall i \in V, \quad n_i = \frac{n_i}{\left\| \overline{n_i} \right\|} \quad average \quad \overline{n_i} = \frac{1}{\left| N(f) \right|} \sum_{f \in N(f)} n_f$$



#### **Vertex Normal**

#### **Max-Nelson's Method**

Take into account the edge lengths

- N(i) number of edges that share the vertex
- (i+1) mod N(i) next edge after i
- Compute the normal vector of vertex i:

$$\forall i \in V, \quad n_i = \frac{1}{|N(f)|} \sum_{j \in N(i)} \frac{e_j \times e_{j+1}}{\|e_j\|^2 \|e_{j+1}\|^2}$$
$$e_j = X_j - X_i$$



#### **Connectivity matrix**

(Discrete Laplacian matrix)

Connectivity graph can be represented as a matrix L with dimension  $N_{\nu} \ge N_{\nu}$ 

$$L_{ij} = \begin{cases} -1 & i = j \\ \lambda_{ij} & (i, j) \in E \text{ (neighbors)} \\ 0 & \text{otherwise} \end{cases}$$

$$L = \begin{pmatrix} -1 & \lambda_{1,2} & \lambda_{1,3} & \lambda_{1,4} & \lambda_{1,5} \\ \lambda_{2,1} & -1 & \lambda_{2,3} & \lambda_{2,4} & 0 \\ \lambda_{3,1} & \lambda_{3,2} & -1 & \lambda_{3,4} & \lambda_{3,5} \\ \lambda_{4,1} & \lambda_{4,2} & \lambda_{4,3} & -1 & 0 \\ \lambda_{5,1} & 0 & \lambda_{5,3} & 0 & -1 \end{pmatrix}$$

choice of weights  $\lambda_{ii}$ 



#### **Connectivity matrix**

(Discrete Laplacian matrix)<sub>1</sub>

$$L_{ij} = \begin{cases} -1 & i = j \\ \lambda_{ij} & (i, j) \in E \text{ (neighbors)} \\ 0 & \text{otherwise} \end{cases}$$

choice of weights  $\lambda_{ij}$  satisfying:

$$\begin{split} \lambda_{i,i} &= -\sum_{j \neq i} \lambda_{i,j}, \\ \lambda_{i,j} &> 0 \quad , \sum_{j \in N(i)} \lambda_{i,j} = 1 \end{split}$$



$$L = \begin{pmatrix} -1 & \lambda_{1,2} & \lambda_{1,3} & \lambda_{1,4} & \lambda_{1,5} \\ \lambda_{2,1} & -1 & \lambda_{2,3} & \lambda_{2,4} & 0 \\ \lambda_{3,1} & \lambda_{3,2} & -1 & \lambda_{3,4} & \lambda_{3,5} \\ \lambda_{4,1} & \lambda_{4,2} & \lambda_{4,3} & -1 & 0 \\ \lambda_{5,1} & 0 & \lambda_{5,3} & 0 & -1 \end{pmatrix}$$



#### **Curvature on surface**

For a given point P each direction  $e_{\Theta}$  in the tangent plane  $T_pS$  defines a curve C (normal section) as the intersection between the plane containing N,  $e_{\Theta}$  and the surface S

Curves passing in different directions have different values of normal curvatures

$$\kappa_N(\theta) = \langle N, C'' \rangle$$

## **Principal curvatures**

For each direction  $e_{\theta} \in T_p S$ , a curve C may have a different normal curvature  $\kappa_N(\theta) = \langle N, C'' \rangle$ 

• A point  $p \in S$  has

multiple curvatures!

Principal curvatures

 $k_1 = \min_{e_{\theta} \in T_p S} \kappa_N(\theta), \quad k_2 = \max_{e_{\theta} \in T_p S} \kappa_N(\theta)$ Principal directions

$$e_1 = \arg\min_{e_{\theta} \in T_p S} \kappa_N(\theta), \quad e_2 = \arg\max_{e_{\theta} \in T_p S} \kappa_N(\theta)$$





K determines if a surface is locally saddle (-) or locally convex (+)



$$H(v) = \frac{1}{2A_{v}} \sum_{j \in N(v)} (\cot \alpha_{j} + \cot \beta_{j}) ||v - v_{j}||$$
$$\vec{H}(v) = \frac{1}{2A_v} \sum_{j \in N(v)} (\cot \alpha_j + \cot \beta_j)(v - v_j)$$

$$\vec{H}(v) = \frac{1}{2A_v} \sum_{j \in N(v)} (\cot \alpha_j + \cot \beta_j)(v - v_j)$$

$$\cot \alpha = \frac{e_1 \cdot e_2}{\sqrt{\|e_1\|^2 \|e_2\|^2 - (e_1 \cdot e_2)^2}}$$

$$e_1 = v - v_{i+1}, \quad e_2 = v_i - v_{i+1}$$
Theorem:  $(\cot(\alpha_{ij}) + \cot(\beta_{ij})) > 0 \quad \leftrightarrow \quad \alpha_{ij} + \beta_{ij} < \pi$ 



### **Discrete Mean Curvature**







#### False colors visualization



Gaussian Curvature at v:

$$K_{v} = \frac{1}{A_{v}} (2\pi - \sum_{j \in N(v)} \theta_{j})$$



where  $\theta_i$  are the incident internal angles at v

- Measure of the distance of a surface (mesh) to a plane
- Depends on angles and lengths
- Zeros curvature in flat areas



### **Principal curvatures**

$$\kappa_1 = H + \sqrt{H^2 - K}$$

$$\kappa_2 = H - \sqrt{H^2 - K}$$



## **Properties of a Mesh**

- **Solidity** : a mesh represents a solid object if its faces together enclose a positive and finite amount of space.
- Connectedness : A mesh is connected if an unbroken path along polygon edges exists between any two vertices.
- **Simplicity** : A mesh is **simple** if the object it represents is solid and has no holes through it ; that is, the object can be deformed into a sphere without tearing. Genus = 0
- Convexity : The mesh represents a convex object if the line connecting any two points within the object lies wholly inside the object.



logna



### **Surface Genus**

In TOPOLOGY, the genus of a surface is defined as the biggest number of simple, close curves that can be drawn on the surface without splitting it into two non-connected parts

For orientable surfaces, the genus counts the number of "handles or holes" of an object







### Euler's formula for a mesh without boundaries

Euler's formula provides a fundamental relationship between the number of faces, edges, and vertices for polyhedral in a closed and connected (but otherwise unstructured) mesh.

For a mesh that is not simple:

$$\chi = |V| - |E| + |T| = 2(1 - g)$$

is the **Euler Characteristic** and **g** is the GENUS through the polyhedron.

For a simple, solid, connected mesh:

$$\chi = |V| - |E| + |T| = 2$$
$$\chi(cube) ??$$



### Euler's formula for mesh with Boundaries

$$\chi = |V| - |E| + |T| + |B| = 2(1 - g)$$

### **|B|** # of boundaries

Example: Genus = 1 - (4 - 5 + 2 + 1)/2= 0



Triangle Meshes with v vertices have about 2v face and 3v edges





- Vertex normals
  - Prefixed w/ "vn" (Wavefront)
  - Contains x,y,z of normal
  - Not necessarily unit length
  - Not necessarily in vertex order
  - Indexed as with vertices
- Texture coordinates
  - Prefixed with "vt" (Wavefront)
  - Not necessarily in vertex order
  - Contains u,v surface parameters
- Faces
  - Uses "/" to separate indices
  - Vertex "/" normal "/" texture
  - Normal and texture optional
  - Can eliminate normal with "//"



Some applications support vertex colors, by putting red, green and blue values after x y and z.

#### Mesh Orientation checking

 If the orientation for an edge is the same for both polys sharing it, then flip the poly outline

- Every poly has counterclockwise outline
- Given by the order of the vertices in a face





v3

v1

OK

F v0 v1 v2

F v1 v3 v2



### Definition

- 1-ring neighborhood N(i)
  - Faces/vertices adjacent at vertex i
  - Valence of a vertex



Number of incident edges at a vertex

#### **Common Mesh Operations**

- Direct Access to the elements
- Ordered Access to the elements From an initial elements walk through the mesh elements by adjacent elements (i.e. along edges, faces, ...)
- *Topological Relations* given a face, find out its edges and its vertices. Given a vertex which is the set of incidents elements.



### **Additional operations**

Edge Split

Adds a vertex to get four triangles

- Edge collapse
   Removes an edge
   Removes a vertex
- Edge Flip

Flip an edge

- beware of valence three



# Winged-Edge data structure

- vertex
  - 1 pointer to one edge
- edge pointers to
  - -2 endpoint vertices
  - -2 faces that share edge
  - 4 edges emanating from its endpoints
- face
  - 1 pointer to one edge





### Halfedge data structure

- vertex
  - 1 pointer to halfedge
- halfedge pointers to
  - 1 incident vertex
  - 1 incident face
  - 1, 2, or 3 halfedges (prec.,next.,opposite)

face

- 1 halfedge (random)





## Geometry Mesh Processing

# Geometry Processing Pipeline



scan



process



### Reconstruction

- Simplification
- Parameterization
- Remeshing
- Smoothing/Fairing
- Deformation/Editing
- Shape Analysis





- Reconstruction
- Simplification
- Parameterization
- Remeshing





- Reconstruction
- Simplification
- Parameterization
- Remeshing
- Smoothing/Fairi
- Deformation/Editing
- Shape Ana





[Botsch and Sorkine]



- Reconstruction
- Simplification
- Parameterization
- Remeshing
- Smoothing/Fairing
- Deformation/Editing
- Shape Analysis
  - Identify/understand important semantic features
  - segmentation, correspondence, symmetry detection, ...





- Reconstruction
- Simplification
- Parameterization
- Remeshing
- Smoothing/Fa
- Deformation/Ec





## **Mesh Optimization**

#### PROBLEM:

the number of required polys to efficiently represent a complex object is huge and redundant.

#### **SOLUTION: optimize the mesh preserving its topology**

- Produce approximations with fewer triangles
  - should be as similar as possible to original
  - want computationally efficient process
- Speed up Rendering
- Less storage
- Easier processing





### **Applications**



### Level of detail hierarchies

Vertices: 81457 Faces:162910







Vertices 60K

#### Vertices:17K

Vertices: 400

### **Applications** Adaptation to hardware capabilities



Kobbelt,Botsch,2008

# What makes a "good" mesh?

- Good approximation of original shape!
- Keep only elements that contribute information about shape
- Add additional information where, e.g., curvature is large





- E.g., all angles close to 60 degrees
- More sophisticated condition: *Delaunay*
- Can help w/ numerical accuracy/stability
- Tradeoffs w/ good geometric approximation\*

\*See Shewchuk, "What is a Good Linear Element"

# What else constitutes a good mesh?

- Another rule: regular vertex degree
- E.g., valence 6 for triangle meshes (equilateral)



• Why? Better polygon shape, important for (e.g.) subdivision:





## **Mesh Optimization**

#### PROBLEM STATEMENT:

Given: 3D model M = (V, E,F)

Point samples V={Pi}, Mesh connectivity F={Tj}

Find: 3D model M' = (V', F') such that

$$- ||M - M'|| < \varepsilon$$

Most common kinds of optimization methods:

- 1) Vertex clustering decimation
- 2) Incremental Decimation



### Incremental mesh Decimation (data reduction)

#### ALGORITHM:

#### Input M (original model)

#### For each geometric item (edge/face)

- rank all geometric item with some cost metric
- sort for increasing cost

#### Repeat

- contract minimum cost geometric item by decimation
- update geometric item list of costs until (no further reduction possible)
   Output M' (simplified model)



### Decimation Operators: vertex remove

### **Remove of the vertex**

- Remove associated triangles
- Triangulate the hole





### Decimation Operators: edge collapse

**Edge collapse**:  $(v1,v2) \implies v$  (based on appropriate rule *like average*) **Vertex split:** removes:

- 1 vertex
- 3 edges
- 2 triangles

Dual of edge collapse, adds 2 triangles







Garland & Heckbert, SIGGRAPH 97 (DirectX)

- Based on point-to-plane distance
- The set of planes at a vertex is initialized to be the planes of the triangles that meet at that vertex.





### **Quadric Error Metric**

• The error of the vertex v with respect to the set of planes is the sum of squared distances from vertex to planes:

$$\Delta_{v} = \sum_{p} Dist(v, p)^{2}$$
$$v = \begin{pmatrix} x \\ y \\ z \\ 1 \end{pmatrix} \quad p = \begin{pmatrix} a \\ b \\ c \\ d \end{pmatrix}$$

Signed distance to plane:

$$Dist(v, p) = ax + by + cz + d = p^T v$$


## **Quadric Error Metric**

	Error metric rewritten as a quadratic form: symmetric 4x4 matrix Q multiplied twice by a vector		
$_{v} = \sum (p^{T} v)^{2}$			
p			
$=\sum v^T p p^T v$	$\int a^2$	ab ac	ad
р	$O^{(p)} - ab$	$b^2$ $bc$	bd
$= v^T \sum Q^{(p)} v$	$\mathcal{Q}$ – ac	$bc c^2$	cd
p T o	ad	bd cd	$d^2$
$= v^{\perp} Q v$			

- Compute Q<sup>(p)</sup> for each triangle (distance to plane p)
- Set Q at each vertex to sum of Qs from incident triangles



## **Using Quadrics**

- Approximate error of edge collapses
  - Each vertex v has associated quadric Q
  - Error of collapsing  $v_1$  and  $v_2$  to  $v_3$  is

$$v_3^T Q_1 v_3 + v_3^T Q_2 v_3$$

– Quadric for new vertex  $v_3$  is  $Q_3=Q_1+Q_2$ 





## **Using Quadrics**

#### Find optimal location $v_3$ after collapse:

 $Q_{3} = \begin{pmatrix} q_{11} & q_{12} & q_{13} & q_{14} \\ q_{12} & q_{11} & q_{23} & q_{24} \\ q_{13} & q_{23} & q_{33} & q_{34} \\ q_{14} & q_{24} & q_{34} & q_{44} \end{pmatrix}$  $\min_{v_{3}} v_{3}^{T} Q_{3} v_{3} : \quad \frac{\partial}{\partial x} = \frac{\partial}{\partial y} = \frac{\partial}{\partial z} = 0$ 

- How do we find a critical point (min/max/saddle)?
- Set derivative to zero!
- matrix Q is *positive-definite*  $\rightarrow$  *min*



## **Using Quadrics**

#### Find optimal location $v_3$ after collapse:

$$Q_{3}v_{3} = \begin{pmatrix} q_{11} & q_{12} & q_{13} & q_{14} \\ q_{12} & q_{11} & q_{23} & q_{24} \\ q_{13} & q_{23} & q_{33} & q_{34} \\ 0 & 0 & 0 & 1 \end{pmatrix} v_{3} = \begin{pmatrix} 0 \\ 0 \\ 0 \\ 0 \\ 1 \end{pmatrix}$$
$$v_{3} = \begin{pmatrix} q_{11} & q_{12} & q_{13} & q_{14} \\ q_{12} & q_{11} & q_{23} & q_{24} \\ q_{13} & q_{23} & q_{33} & q_{34} \\ 0 & 0 & 0 & 1 \end{pmatrix}^{-1} \begin{pmatrix} 0 \\ 0 \\ 0 \\ 1 \end{pmatrix}$$



## Incremental mesh Decimation (with Quadric Error Metric)

ALGORITHM: INPUT: the original model

• Compute the **Q** matrices for all the initial vertices

#### For each edge

Compute the optimal contraction target v for each edge (v1,v2). The error v<sup>T</sup>(Q1+Q2) v of this target vertex becomes the cost of collapsing that edge.

Sort for increasing cost

#### Repeat

- Edge collapse the minimum cost edge (vi,vj) (decimation) to get new vertex v
- add Qi and Qj to get quadric Qv at v
- update cost of edges touching v
- until (no further reduction possible)

# **Results:** Quadric Visualization



#### Progressive Mesh Algorithm (Hoppe 1996)









- Store edge collapses in a ordered list (decreasing edge collapsing (ecol) cost)
- Iteratively process a decimation of minimum cost and recompute the list of costs for neighbors
- An edge collapse is ok only if it does not change the mesh topology





# Mesh simplification applied to .. LOD Approximation Progressive Transmission Mesh compression Selective Refinement



## **LOD Approximation**



#### Create meshes at different Level of Detail (LOD)

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## **LOD Approximation**

#### **Need Multiresolution Models**

- Context dictates required detail
  - LOD should vary with context



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#### need high detail near the viewer need less detail far away



Example: obj with 10000T viewer near the object, if the object is far away (covers10x5 pixels) use a simplified model with 100 T



## LOD

#### • LOD of parametric surfaces:

Bézier patches, spline and subdivision have a natural LOD on discrete parametric domain.Different LODs are given by the resolution of the tessellation

LOD of meshes:

ALGORITHM:

- Generation: Given a model, build a set of approximations can be produced by any simplification system
- Selection: at run time, simply select which to render
- Switching: Inter-frame switching from a LOD to the next



## LOD switching

- Discrete geometry LOD
  - Switch from a LOD at frame i to the next LOD at frame i+1
  - Causes popping
- Blend LOD
  - Transition from LOD 1 to LOD2 in FB. If LOD1 is the current LOD then draw LOD1 in z-buffer and color buffer with alpha =1 (opacity)
  - Draw LOD2 in FB with alpha =0 (trasparency)
  - LOD2 appears incrementing alpha from 0 to 1
  - When LOD2 is visualized (alpha =1) LOD1 start to disappear (alpha decreases to 0)
  - Inbetween both LODs are rendered overlapped
  - No popping, it works in hardware





## Evaluate a criterium based on point of view and object position, choose a suitable LOD

#### Range-based

- Distance from point of view
- Each LOD is associated with a distance range (LOD0 more detailed associated with [0,r1])



## **LOD** selection

#### Projected area based

- -Projected Area from bounding volume of the object (eg.sphere)
- -Sphere center c radius r, camera at v, direction vector d,
- -n distance camera-near plane





## **3D Data (Mesh) Compression**

#### Different needs:

Lossy...



Games



Virtual malls

o Lossless?





## **Mesh compression**



#### Not only the number of polys is reduced, also the storage for each LOD model is compressed.



The user waits for the entire transmission time in order to visualize the entire object,

Single-rate...



#### <u>ransmission</u>

The user sees a first grasp of the object very early during the transmission.



Coarse mesh +Database edge-collapse (compresses)

**APPLY REVERSIBLE DECIMATIONS** 



### **Selective Refinement**

# We may need varying LOD over surface:

- large surface, oblique view (eg. on terrain)
  - high detail near the viewer
  - less detail far away

#### single LOD will be

inappropriate

- either excessively detailed in the distance (wasteful)
- or insufficiently detailed near viewer (visual artifacts)





# Selective Refinement for territorial models

- Find a base mesh
- Subdivide recursively the base mesh to approximate the original mesh







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## Outline

- Reconstruction
- Simplification
- Parameterization
- Remeshing
- Smoothing/Fairing
- Deformation/Editing
- Shape Analysis





#### **Parameterization**

A parameterization is a bijective mapping between a surface and a parameter domain



- surface  $S \subset \mathbb{R}^3$
- parameter domain  $\Omega \subset \mathbb{R}^2$
- mapping  $f: \Omega \to S$  and  $f^{-1}: S \to \Omega$



### **Mesh Parameterization**







## **Applications - Remeshing**



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#### **Desirable Properties**



- Bijective function f (1-1 and onto): No triangles fold over.
- Minimal "distortion"
  - Preserve 3D angles (conformal)
  - Preserve 3D lengths (isometric)
  - Preserve 3D areas (equiareal)
  - No "stretch"
- Efficiently computable





- f is isometric (length preserving), if the length of any arc on S is the same as that of its pre-image on  $\Omega$ .
- f is conformal (angle preserving), if the angle of intersection of every pair of intersecting arcs on S\* is the same as that of the corresponding preimages on Ω.
- f is equiareal (area preserving) if every part of Ω
   is mapped onto a part of S\* with the same area

Distortion is (almost) Inevitable

• Theorema Egregium (C. F. Gauß)

"A general surface cannot be parameterized without distortion."

- no distortion = conformal + equiareal = isometric
- requires surface to be developable
  - planes
  - cones
  - cylinders



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### **Genus=0**, with boundary







Introduce cuts on the mesh





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#### **UNFOLDING the word**





Mario Botsch, Bielefeld Graphics & Geometry ALMA MATER STUDIORUM - UNIVERSITÀ DI BOLOGNA



#### Convex combination maps (Tutte-Floater)

#### **Barycentric mapping theorem (Tutte)**

Given a triangulated surface homeomorphic to a disk, if the (u,v) coordinates at the boundary vertices lie on a convex polygon, and if the coordinates of the internal vertices are a convex combination of their neighbors, then the (u,v) coordinates form a valid parametrization (without self-intersection).



ALMA MATER STUDIORUM - UNIVERSITÀ DI BOLOGNA Eck, DeRose, Duchamp, Hoppe, Lounsbery, Stuetzle: Multiresolution Analysis of Arbitrary Meshes, SIGGRAPH, 1995

#### Triangle Mesh Parameterization Convex combination maps (Tutte-Floater)

- Works for meshes equivalent to a disk
- Algorithm:
  - First, we map the 3D boundary to a 2D convex polygon
  - Then we compute the inner vertices positions



# **Compute the inner vertices**

• We constrain each inner vertex to be a weighted average of its neighbors:

$$x_i = \sum_{j \in N(i)} \lambda_{i,j} x_j, \quad i = 1, 2, ..., n$$

with weights

 $\begin{cases} 0 & \text{if } (i, j) \text{ are not neighbors} \\ \lambda_{i,j} > 0 & \text{if } (i, j) \in E \text{ (neighbors)} \\ \lambda_{i,i} = -\sum_{j \neq i} \lambda_{i,j}, \\ \sum_{j \in N(i)} \lambda_{i,j} = 1 \end{cases}$ 





$$x_{i} - \sum_{j \in N(i)} \lambda_{i,j} x_{j} = 0, \quad i = 1, 2, ..., n$$

$$x_{i} - \sum_{j \in N(i) \setminus B} \lambda_{i,j} x_{j} = \sum_{\substack{k \in N(i) \cap B \\ \text{fixed } (r_{i})}} \lambda_{i,k} x_{k}, \quad i = 1, 2, ..., n$$

$$\underbrace{\left( \begin{array}{c} -1 & \lambda_{1,j_{1}} & \lambda_{1,j_{d_{1}}} \\ -1 & & \\ \lambda_{4,j_{1}} & \ddots & \\ & & \\ \end{array} \right)}_{L} \left( \begin{array}{c} x_{1} \\ x_{2} \\ \vdots \\ x_{n} \end{array} \right) = \begin{pmatrix} -r_{1} \\ -r_{2} \\ \vdots \\ x_{n} \end{pmatrix} = \begin{pmatrix} -r_{1} \\ -r_{2} \\ \vdots \\ x_{n} \end{pmatrix} = \begin{pmatrix} -r_{1} \\ -r_{2} \\ \vdots \\ x_{n} \end{pmatrix} = \begin{pmatrix} -r_{1} \\ -r_{2} \\ \vdots \\ x_{n} \end{pmatrix} = \begin{pmatrix} -r_{1} \\ -r_{2} \\ \vdots \\ x_{n} \end{pmatrix} = \begin{pmatrix} -r_{1} \\ -r_{2} \\ \vdots \\ x_{n} \end{pmatrix} = \begin{pmatrix} -r_{1} \\ -r_{2} \\ \vdots \\ x_{n} \end{pmatrix} = \begin{pmatrix} -r_{1} \\ -r_{2} \\ \vdots \\ x_{n} \end{pmatrix} = \begin{pmatrix} -r_{1} \\ -r_{2} \\ \vdots \\ x_{n} \end{pmatrix} = \begin{pmatrix} -r_{1} \\ -r_{2} \\ \vdots \\ x_{n} \end{pmatrix} = \begin{pmatrix} -r_{1} \\ -r_{2} \\ \vdots \\ x_{n} \end{pmatrix} = \begin{pmatrix} -r_{1} \\ -r_{2} \\ \vdots \\ x_{n} \end{pmatrix} = \begin{pmatrix} -r_{1} \\ -r_{2} \\ \vdots \\ x_{n} \end{pmatrix} = \begin{pmatrix} -r_{1} \\ -r_{2} \\ \vdots \\ x_{n} \end{pmatrix} = \begin{pmatrix} -r_{1} \\ -r_{2} \\ \vdots \\ -r_{n} \end{pmatrix}$$


## Linear system of equations

• solve system twice

$$x_i = (u_i, v_i) \Longrightarrow$$

• matrix L is

$$LU = U_B$$
 and  $LV = V_B$ 

- sparse
- diagonally dominant
- nonsingular
- A unique solution always exists
- Important: the solution is legal (bijective)
- The system is sparse, thus fast numerical solution is possible
- Numerical problems (because the vertices in the middle might get very dense...)



## **Choice of Weights**

 $A_i$ 

#### Discrete Laplacian

 $L = (\lambda_{i,j}), \qquad \lambda_{i,i} = -1$ 

- The weights are given by
- Uniform /Barycentric (Tutte)

 $\alpha_{ii}$ 

 $\overline{\lambda_{i,j}} = \frac{1}{d_i}$  d<sub>i</sub> valence of vertex i

- Discrete
  harmonic
- Mean value

 $\overline{\lambda_{i,j}} = \frac{1}{2A_i} (\cot \alpha_{i,j} + \cot \beta_{i,j})$   $\overline{\lambda_{i,j}} = \frac{1}{\|X_i - X_j\|} (\tan(\frac{\delta_{i,j}}{2}) + \tan(\frac{\gamma_{i,j}}{2}))$   $\lambda_{i,j} = \frac{\overline{\lambda_{i,j}}}{\sum_{k \in N(i)} \overline{\lambda_{i,k}}}$ SITÀ DI BOLOGNA



## **Convex Combination Maps**

Comparison









(S. Yoshizawa , 2004)

Iterative algorithm to correct the weights

- Initialize: convex map
- STEP *k*:
  - Update  $\lambda_{i,j}$  w.r.t. stretch  $\sigma_j$  in  $X_j$

$$\lambda_{i,j}^k = \frac{\overline{\lambda_{i,j}}}{\sum_{n \in N(i)} \overline{\lambda_{i,n}}} \quad \overline{\lambda_{i,j}} = \frac{\lambda_{i,j}^{k-1}}{\sigma_j}$$

Solve the new LS

 Stopping when the global stretch energy stops decreases



## S D JO RUM

## **Texture Atlas Generation**

Split model into number of patches (atlas)

- because higher genus models cannot be mapped onto plane and/or
- because distortion will be too high eventually



Levy, Petitjean, Ray, Maillot: Least Squares Conformal Maps for Automatic Texture Atlas Generation, SIGGRAPH, 2002





Levy: Constraint Texture Mapping, SIGGRAPH, 2001



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## **Surface fairing**



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## **Mesh Fairing**



**X** is a surface parameterization of **S** 

X is the vertex set

of M



Fairness: low variation in curvature.

Move vertices to achieve - Mesh topology stays the same.





- A surface smoothing method, named fairing, removes undesirable noise and uneven edges from discrete surfaces.
- Extend the low-pass filter in image processing
- Idea: a low pass filter corresponds to diffusion flow.



## Image Denoising by linear diffusion flow

$$\frac{\partial u}{\partial t} = \Delta u \quad u_0 = \text{noisy image}$$

$$u(t) = G_{\sqrt{2t}} * u(0) \qquad \sigma = \sqrt{2t}$$



[heat equation = Gaussian filter Witkin-Koenderink '83-84] PLOGNA



## **Diffusion Flow on Meshes**

#### 1) Mean Curvature Flow (MCF)



#### 2) Surface Diffusion Flow

$$\frac{\partial X}{\partial t} = \Delta_s H(\mathbf{X})$$

3) MCF + data fidelity

$$\frac{\partial X}{\partial t} = \Delta_S X + \lambda (f - X^0), \quad X \big|_{t=0} = X^0$$



# Noisy mesh MCF

### SDF

and its curvature



igea 268K faces









## **Discrete diffusion flow on mesh**

- Considering a uniform discretization of the time interval [0, T], T > 0, and using a temporal time step *τ*,
- the approximation of an evolving surface at the n-th time step is denoted by a spatial position vector X<sup>n</sup>

$$\frac{\partial X}{\partial t} \approx \frac{X^{n+1} - X^n}{\tau}$$

•  $L \in \mathbb{R}^{N_V \times N_V}$  (connectivity matrix) represents the discrete Laplace-Beltrami operator

$$\Delta_M X \approx L X$$

In matrix vector form



## **Discrete diffusion flow on mesh**

• Considering a uniform discretization of the time interval [0, T], T > 0, and using a temporal time step  $\tau$ , the approximation of an evolving surface at the n-th time step by MCF  $\partial X$ 

$$\frac{\partial X}{\partial t} = \Delta_s X$$

- Given the position vector  $X^n$  we get  $X^{n+1}$  by applying explicit schemes:  $X^{n+1} = X^n + \tau L^n X^n$
- or implicit schemes  $(I \tau L^n)X^{n+1} = X^n$

where L represents the discrete Laplace-Beltrami operator



Laplace Beltrami operator applied to

the coordinate function X = (x,y,z)

Desbrun et al., Siggraph 1999 Desbrun et al., Siggraph 1999



Note that since the Laplacian measures the difference between a vertex and the average values of its neighbors, we expect that is equal to zero on constant functions c, so:

 $\Rightarrow Lc = 0$ 



Results of applying **umbrella** weights on the left and **cotangent** weights on the right.



## Material

#### Level Set Methods and Dynamic Implicit Surfaces

Stanley Osher & Ronald Fedkiw Series: Applied Mathematical Sciences, Vol. 153 Springer, 2003 ISBN-10: 0-387-95482-1 ISBN-13: 978-0-387-95482-0



#### **Polygon Mesh Processing**

Mario Botsch, Leif Kobbelt Mark Pauly, Pierre Alliez, Bruno Levy A K Peters, Ltd. Natick, Massachusetts 2008 ISBN 978-1-56881-426-1







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