

Bounds for the Matrix Condition Number

**CIME-EMS School, Exploiting Hidden Structure
in Matrix Computations. Algorithms and Applications**

Cetraro 22-26 June 2015

Sarah W. Gaaf
Joint work with Michiel Hochstenbach (TU/e)

Uses of $\kappa(A)$

Consider $A\mathbf{x} = \mathbf{b}$, $A \in \mathbb{R}^{n \times n}$, A nonsingular.

Problem:

- A , \mathbf{b} perturbed
- Compute $\mathbf{x} = A^{-1}\mathbf{b}$
- Accuracy solution \mathbf{x} ?

Sensitivity of linear system:

$$\frac{\|\Delta\mathbf{x}\|}{\|\mathbf{x}\|} \leq \kappa(A) \left(\frac{\|\Delta A\|}{\|A\|} + \frac{\|\Delta\mathbf{b}\|}{\|\mathbf{b}\|} \right)$$

How to compute the condition number?

Matrix 2-norm:

$$\|A\|_2 = \max_{x \neq 0} \frac{\|Ax\|_2}{\|x\|_2} = \sigma_1(A) = \sqrt{\lambda_1(A^T A)} = \lambda_1 \left(\begin{bmatrix} 0 & A^T \\ A & 0 \end{bmatrix} \right)$$

Matrix condition number:

$$\kappa(A) = \|A\|_2 \|A^{-1}\|_2 = \frac{\sigma_1(A)}{\sigma_n(A)} = \sqrt{\frac{\lambda_1(A^T A)}{\lambda_n(A^T A)}}$$

Bounds for largest singular value

Lower bound θ :

- Krylov subspace methods (Lanczos bidiagonalization)

Upper bounds:

- Bauer-Fike theorem (1960): *There is a singular value in $[\theta, \theta + C]$.*
 - not guaranteed σ_1
- Probabilistic bound [Hochstenbach 2013]

Problem:

- Find upper and lower bound for smallest singular value

Extended Lanczos Bidiagonalization

Procedure:

$$\mathbf{v}_0 \xrightarrow{A} \mathbf{u}_0 \xrightarrow{A^T} \mathbf{v}_1 \xrightarrow{A^{-T}} \mathbf{u}_1 \xrightarrow{A^{-1}} \dots$$

Equations:

$$\begin{aligned} A^T A V &= V H^T H \\ (A^T A)^{-1} V &= V K K^T & \mathbf{v}_k &= p_k(A^T A) \mathbf{v}_0 \\ A A^T U &= U H H^T & \mathbf{u}_k &= q_k(A A^T) \mathbf{u}_0 \\ (A A^T)^{-1} U &= U K^T K \end{aligned}$$

Characteristics:

- H and K are tridiagonal, $H \cdot K = I$
- p_k and q_k are Laurent polynomials
- $\mathcal{K}_{m,m}(A^T A, \mathbf{v}_0) = \text{span}\{\dots, (A^T A)^{-1} \mathbf{v}_0, \mathbf{v}_0, A^T A \mathbf{v}_0, \dots\}$.

Lower bound for $\kappa(A)$

Extended Lanczos Bidiagonalization:

- Largest singular value θ_1 of H approximates $\sigma_1(A)$
- Smallest singular value θ_k of H approximates $\sigma_n(A)$
- Lower bound for condition number:

$$\frac{\theta_1}{\theta_k} \leq \kappa(A)$$

Probabilistic upper bound

Let $\mathbf{v}_0 = \sum_{i=1}^n \gamma_i \mathbf{y}_i$, (\mathbf{y}_i right singular vectors of A)

then

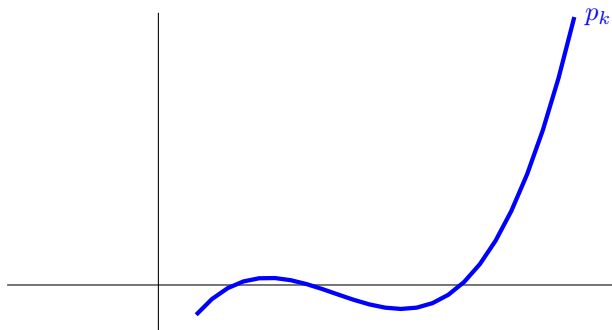
$$1 = \|\mathbf{v}_k\|^2 = \|p_k(A^T A)\mathbf{v}_0\|^2 = \sum_{i=1}^n \gamma_i^2 p_k(\sigma_i^2)^2.$$

Thus

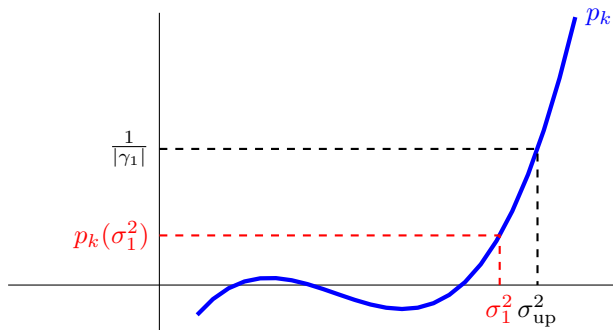
$$1 \geq \gamma_1^2 p_k(\sigma_1^2)^2,$$

and

$$\frac{1}{|\gamma_1|} \geq |p_k(\sigma_1^2)|.$$



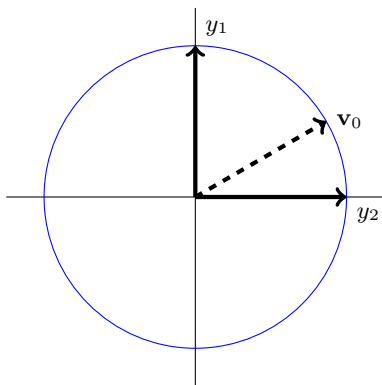
Recall: $\frac{1}{|\gamma_1|} \geq |p_k(\sigma_1^2)|$



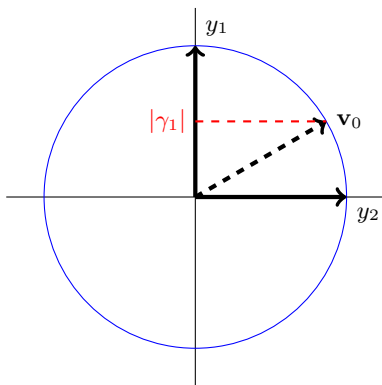
Recall: $\frac{1}{|\gamma_1|} \geq |p_k(\sigma_1^2)|$

Question: $\mathbb{P}(\frac{1}{\delta} < \frac{1}{|\gamma_1|}) = \mathbb{P}(|\gamma_1| < \delta) = \epsilon$ (ϵ is user-chosen)

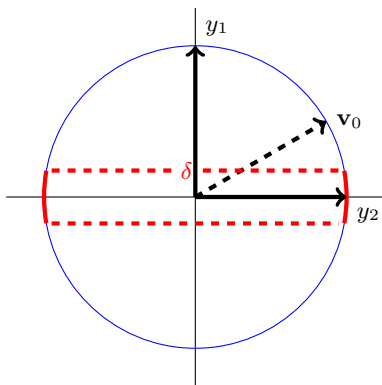
$$\mathbb{P}\left(\frac{1}{\delta} < \frac{1}{|\gamma_1|}\right) = \mathbb{P}(|\gamma_1| < \delta) = \epsilon \quad (\epsilon \text{ is user-chosen})$$



$$\mathbb{P}\left(\frac{1}{\delta} < \frac{1}{|\gamma_1|}\right) = \mathbb{P}(|\gamma_1| < \delta) = \epsilon \quad (\epsilon \text{ is user-chosen})$$



$$\mathbb{P}\left(\frac{1}{\delta} < \frac{1}{|\gamma_1|}\right) = \mathbb{P}(|\gamma_1| < \delta) = \epsilon \quad (\epsilon \text{ is user-chosen})$$



Theorem:

- starting vector \mathbf{v}_0 chosen randomly (uniform distribution over S^{n-1})
- $\varepsilon \in (0, 1)$ user chosen
- δ be given by $\varepsilon = \mathcal{P}(|\gamma_1| \leq \delta) = \frac{B_{\text{inc}}(\frac{n-1}{2}, \frac{1}{2}, \delta^2)}{B_{\text{inc}}(\frac{n-1}{2}, \frac{1}{2}, 1)}$,

where $B_{\text{inc}}(x, y, z) = \int_0^z t^{x-1} (1-t)^{y-1} dt$ (incomplete Beta function)

Then $\sigma_{\text{up}}^{\text{prob}}$, the square root of the largest zero of the polynomial

$$f_1^\delta(t) = |p_k(t)| - \frac{1}{\delta},$$

is upper bound for σ_1 with probability at least $1 - \varepsilon$.

Conclusions

Bounds for the condition number:

- lower bound: $\kappa_{\text{low}} = \frac{\theta_1}{\theta_k} \leq \kappa(A)$
- probabilistic upper bound: $\kappa(A) \leq \frac{\sigma_1^{\text{prob}}}{\sigma_n^{\text{prob}}} = \kappa_{\text{up}}$

User chosen values:

- ε , probabilistic bound holds with probability at least $1 - 2\varepsilon$
- ζ , method adaptively performs k steps such that

$$\frac{\kappa_{\text{up}}}{\kappa_{\text{low}}} \leq \zeta$$

[SG, M. E. Hochstenbach, *Probabilistic bounds for the matrix condition number with extended Lanczos bidiagonalization* (Submitted)]

Matrix A	Dim.	κ	κ_{low}	κ_{up}	k	CPU	LU	CPU ¹
utm5940	5940	$4.35 \cdot 10^8$	$3.98 \cdot 10^8$	$7.21 \cdot 10^8$	4	0.13	61	0.12
grcar10000	10000	$3.63 \cdot 10^0$	$3.59 \cdot 10^0$	$5.80 \cdot 10^0$	6	0.07	31	0.05
af23560	23560	$1.99 \cdot 10^4$	$1.93 \cdot 10^4$	$2.82 \cdot 10^4$	6	0.98	74	0.88
rajat16	96294	*	$5.63 \cdot 10^{12}$	$5.69 \cdot 10^{12}$	5	9.34	97	9.19
torso1	116158	*	$1.41 \cdot 10^{10}$	$1.42 \cdot 10^{10}$	3	26.8	93	28.5
dc1	116835	*	$2.39 \cdot 10^8$	$4.59 \cdot 10^8$	5	6.05	93	5.57
xenon2	157464	*	$4.29 \cdot 10^4$	$8.14 \cdot 10^4$	7	20.1	82	19.6
scircuit	170998	*	$2.40 \cdot 10^9$	$4.69 \cdot 10^9$	7	2.05	54	1.39
transient	178866	*	$1.02 \cdot 10^{11}$	$2.00 \cdot 10^{11}$	8	7.70	86	7.12
stomach	213360	*	$4.62 \cdot 10^1$	$9.02 \cdot 10^1$	6	13.8	80	13.7

- For $\zeta = 2$ (i.e. $\kappa_{\text{up}}/\kappa_{\text{low}} \leq 2$)
- $\varepsilon = 0.01$ (i.e. upper bound holds with probability at least 98%)
- CPU¹ indicates time of `condst`

Matrix A	Dim.	κ	κ_{low}	κ_{up}	k	CPU	LU
utm5940	5940	$4.35 \cdot 10^8$	$4.35 \cdot 10^8$	$4.71 \cdot 10^8$	10	0.19	42
grcar10000	10000	$3.63 \cdot 10^0$	$3.62 \cdot 10^0$	$3.97 \cdot 10^0$	13	0.13	21
af23560	23560	$1.99 \cdot 10^4$	$1.99 \cdot 10^4$	$2.12 \cdot 10^4$	9	1.14	66
rajat16	96294	*	$5.63 \cdot 10^{12}$	$5.69 \cdot 10^{12}$	5	9.34	97
torso1	116158	*	$1.41 \cdot 10^{10}$	$1.42 \cdot 10^{10}$	3	26.8	93
dc1	116835	*	$2.39 \cdot 10^8$	$2.45 \cdot 10^8$	8	6.52	91
xenon2	157464	*	$4.32 \cdot 10^4$	$4.67 \cdot 10^4$	14	23.6	70
scircuit	170998	*	$2.45 \cdot 10^9$	$2.67 \cdot 10^9$	16	3.28	33
transient	178866	*	$1.03 \cdot 10^{11}$	$1.11 \cdot 10^{11}$	21	9.47	70
stomach	213360	*	$4.82 \cdot 10^1$	$5.24 \cdot 10^1$	14	17.5	63

- For $\zeta = 1.1$ (i.e. $\kappa_{\text{up}}/\kappa_{\text{low}} \leq 1.1$)
- $\varepsilon = 0.01$ (i.e. upper bound holds with probability at least 98%)

References

HOCHSTENBACH, 2013

HALKO, MARTINSSON, AND TROPP, 2011

JAGELS AND REICHEL, 2011

KNIZHNERMAN AND SIMONCINI, 2010

JAGELS AND REICHEL, 2009

SIMONCINI, 2007

VAN DORSSELAER, HOCHSTENBACH, AND VAN DER VORST, 2000

DRUSKIN AND KNIZHNERMAN, 1998

KUCZYŃSKI AND WOŹNIAKOWSKI, 1992