

Lecture 1. Matrix-Tensor Connections

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**Using examples, let us first take
a look at what we might
mean by
“hidden structure”
in a matrix.**

Hidden Matrix Structure: Five Motivating Examples

The Discrete Fourier Transform

Definition

$$F_4 = \begin{bmatrix} 1 & 1 & 1 & 1 \\ 1 & \omega_4 & \omega_4^2 & \omega_4^3 \\ 1 & \omega_4^2 & \omega_4^4 & \omega_4^6 \\ 1 & \omega_4^3 & \omega_4^6 & \omega_4^9 \end{bmatrix} \quad \omega_n = \cos\left(\frac{2\pi}{n}\right) - i \sin\left(\frac{2\pi}{n}\right)$$

Hidden Structure

$$F_{2m} \Pi_{2,m} = \begin{bmatrix} F_m & \Omega_m F_m \\ F_m & -\Omega_m F_m \end{bmatrix} \quad \begin{array}{l} \Pi_{2,m} = \text{perfect shuffle} \\ \Omega_m = \text{diagonal} \end{array}$$

Recursive Block Structure

Hidden Matrix Structure: Five Motivating Examples

The DFT Matrix is Data Sparse

The DFT matrix is dense, but can be factored into a product of sparse matrices:

$$F_{1024} = A_{10} \cdots A_2 A_1 P^T$$

The A_k have the form $I \otimes \begin{bmatrix} I & D \\ I & -D \end{bmatrix}$, $D = \text{diagonal}$.

That is what makes the FFT possible:

$$\begin{aligned} y &= x \\ \text{for } k &= 1:10 \\ y &= A_k y \end{aligned}$$

An N -by- N matrix is data sparse if it can be represented with many fewer than N^2 numbers. F_N is data sparse: $O(N \log N)$ vs $O(N^2)$.

2. Hamiltonian Matrices

Definition

$$M = \begin{bmatrix} A & F \\ G & -A^T \end{bmatrix} \quad F = F^T, G = G^T$$

Hidden Structure

$$M \begin{bmatrix} y \\ z \end{bmatrix} = \lambda \begin{bmatrix} y \\ z \end{bmatrix} \Rightarrow M^T \begin{bmatrix} z \\ -y \end{bmatrix} = -\lambda \begin{bmatrix} z \\ -y \end{bmatrix}$$

Eigenvalues come in plus-minus pairs.

Hidden Matrix Structure: Five Motivating Examples

Transformations that Preserve Structure

Equivalent Definition

$$J_{2n}^T M J_{2n} = -M^T \quad J_{2n} = \begin{bmatrix} 0 & I_n \\ -I_n & 0 \end{bmatrix}$$

Structured Schur Decomposition

$$\begin{bmatrix} Q_1 & Q_2 \\ -Q_2 & Q_1 \end{bmatrix}^T M \begin{bmatrix} Q_1 & Q_2 \\ -Q_2 & Q_1 \end{bmatrix} = \begin{bmatrix} T_{11} & T_{12} \\ 0 & -T_{11}^T \end{bmatrix}$$

Q is orthogonal and symplectic ($J_{2n}^T Q J_{2n} = Q^{-T}$)

Consequence: Efficient methods for Riccati equations and various “nearness” problems.

3. Cauchy Matrices

Definition

$$A = (a_{kj}) = \left(\frac{1}{\omega_k - \lambda_j} \right) = \begin{bmatrix} \frac{1}{\omega_1 - \lambda_1} & \frac{1}{\omega_1 - \lambda_2} & \frac{1}{\omega_1 - \lambda_3} & \frac{1}{\omega_1 - \lambda_4} \\ \frac{1}{\omega_2 - \lambda_1} & \frac{1}{\omega_2 - \lambda_2} & \frac{1}{\omega_2 - \lambda_3} & \frac{1}{\omega_2 - \lambda_4} \\ \frac{1}{\omega_3 - \lambda_1} & \frac{1}{\omega_3 - \lambda_2} & \frac{1}{\omega_3 - \lambda_3} & \frac{1}{\omega_3 - \lambda_4} \\ \frac{1}{\omega_4 - \lambda_1} & \frac{1}{\omega_4 - \lambda_2} & \frac{1}{\omega_4 - \lambda_3} & \frac{1}{\omega_4 - \lambda_4} \end{bmatrix}$$

Hidden Structure

$$\Omega A - A \Lambda = \text{Rank-1} \quad \Omega = \text{diag}(\omega_i), \Lambda = \text{diag}(\lambda_i),$$

With respect to Ω and Λ , A has displacement rank equal to one.

Hidden Matrix Structure: Five Motivating Examples

Fast LU

First Step:

$$A = \begin{bmatrix} 1 & 0 & 0 & 0 \\ \ell_{21} & 1 & 0 & 0 \\ \ell_{31} & 0 & 1 & 0 \\ \ell_{41} & 0 & 0 & 1 \end{bmatrix} \begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & b_{22} & b_{23} & b_{24} \\ 0 & b_{32} & b_{33} & b_{34} \\ 0 & b_{42} & b_{43} & b_{44} \end{bmatrix} \begin{bmatrix} u_{11} & u_{12} & u_{13} & u_{14} \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \end{bmatrix}$$

Do not compute B as the usual rank-1 update of $A(2:n, 2:n)$. That would be $O(n^2)$

Instead, use the fact that B has unit displacement rank.

The displacement rank representation of B costs $O(n)$

By working with a clever representations it is sometimes possible to dramatically improve efficiency.

Hidden Matrix Structure: Five Motivating Examples

4. Matrices with Orthonormal Columns

Definition

$$Q = \begin{bmatrix} Q_1 \\ Q_2 \end{bmatrix} \quad Q_1^T Q_1 + Q_2^T Q_2 = I$$

Hidden Structure

$$\begin{bmatrix} U_1 & 0 \\ 0 & U_2 \end{bmatrix}^T \begin{bmatrix} Q_1 \\ Q_2 \end{bmatrix} V = \begin{bmatrix} \text{diag}(c_i) \\ \text{diag}(s_i) \end{bmatrix} \quad c_i^2 + s_i^2 = 1$$

$$U_1, U_2, V = \text{orthogonal}$$

Q_1 and Q_2 have related SVDs. This is the CS Decomposition.

Hidden Matrix Structure: Five Motivating Examples

Simultaneous Diagonalization of A_1 and A_2

1. QR factorization: $\begin{bmatrix} A_1 \\ A_2 \end{bmatrix} = \begin{bmatrix} Q_1 \\ Q_2 \end{bmatrix} R$

2. CS decomposition: $\begin{bmatrix} U_1 & 0 \\ 0 & U_2 \end{bmatrix}^T \begin{bmatrix} Q_1 \\ Q_2 \end{bmatrix} V = \begin{bmatrix} \text{diag}(c_i) \\ \text{diag}(s_i) \end{bmatrix}$

3. Setting $X = R^T V$ gives the generalized singular value decomposition:

$$A_1 = U_1 \cdot \text{diag}(c_i) \cdot X^T \quad A_2 = U_2 \cdot \text{diag}(s_i) \cdot X^T$$

An example where exploiting the hidden structure of Q_1 and Q_2 ensures numerical stability.

5. Block Matrices

Definition

$$A = \begin{bmatrix} A_{11} & A_{12} & \cdots & A_{1N} \\ A_{21} & A_{22} & \cdots & A_{2N} \\ \vdots & \vdots & \ddots & \vdots \\ A_{M1} & A_{M2} & \cdots & A_{MN} \end{bmatrix}. \quad \text{Each } A_{ij} \text{ a matrix.}$$

Hidden Structure

The data in each A_{ij} is contiguous in memory.

Not a hidden “math” structure but a “man-made” hidden data structure.

Hidden Matrix Structure: Five Motivating Examples

Respect Data Layout to Minimize Memory Traffic

$$A \leftarrow \begin{bmatrix} A_{11}^T & A_{12}^T & \cdots & A_{1N}^T \\ A_{21}^T & A_{22}^T & \cdots & A_{2N}^T \\ \vdots & \vdots & \ddots & \vdots \\ A_{M1}^T & A_{M2}^T & \cdots & A_{MN}^T \end{bmatrix} . \quad \text{Overwrite } A_{ij} \text{ with } A_{ij}^T .$$

$$A \leftarrow \begin{bmatrix} A_{11}^T & A_{21}^T & \cdots & A_{M1}^T \\ A_{12}^T & A_{22}^T & \cdots & A_{M2}^T \\ \vdots & \vdots & \ddots & \vdots \\ A_{1N}^T & A_{2N}^T & \cdots & A_{MN}^T \end{bmatrix} . \quad \text{Swap } A_{ij}^T \text{ with } A_{ji}^T$$

A 2-pass transpose that exploits the “hidden” data structure.

Each of these examples has a connection to our agenda:

Monday

Lecture 1. Matrix-tensor Connections

Lecture 2. Tensor Symmetries and Rank

Tuesday

Lecture 3. The Tucker and Tensor Train Representations

Lecture 4. The CP and KSVD Representations

Thursday

Lecture 5. Unfolding a Tensor with Multiple Symmetries

Lecture 6. A Higher-Order GSVD

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Data Sparsity

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Structured Permutation Similarity

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A Higher-Order CS Decompositions

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Blocking for Data Locality

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Clever Representations

Let us Begin!

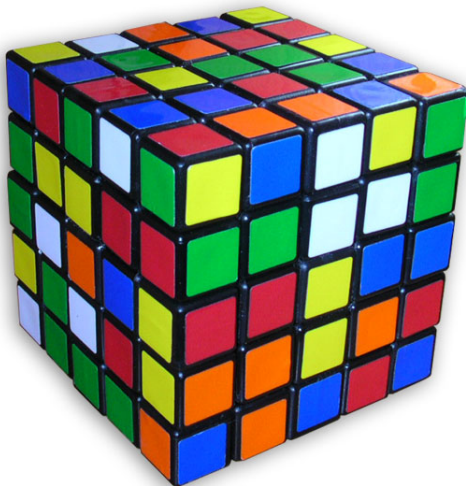
Matrix-Tensor Connections

$$U^T \quad \text{[Rubik's Cube]} \quad V =$$

8	3	5	4	1	6	9	2	7
2	9	6	8	5	7	4	3	1
4	1	7	2	9	3	6	5	8
5	6	9	1	3	4	7	8	2
1	2	3	6	7	8	5	4	9
7	4	8	5	2	9	1	6	3
6	5	2	7	8	1	3	9	4
9	8	1	3	4	5	2	7	6
3	7	4	9	6	2	8	1	5

Much of the discussion will revolve around SVD-like operations.

What is a Tensor?



What is a Tensor?

Definition

An order- d tensor $\mathcal{A} \in \mathbb{R}^{n_1 \times \dots \times n_d}$ is a real d -dimensional array

$$\mathcal{A}(1:n_1, \dots, 1:n_d)$$

where the index range in the k -th **mode** is from 1 to n_k .

Low-Order Tensors

A scalar is an order-0 tensor.

A vector is an order-1 tensor.

A matrix is an order-2 tensor.

We use calligraphic font to designate tensors that have order 3 or greater
e.g., \mathcal{A} , \mathcal{B} , \mathcal{C} , etc.

Fibers

A fiber of a tensor \mathcal{A} is a vector obtained by fixing all but one \mathcal{A} 's indices. For example, if $\mathcal{A} = \mathcal{A}(1:3, 1:5, 1:4, 1:7)$, then

$$\mathcal{A}(2, :, 4, 6) = \mathcal{A}(2, 1:5, 4, 6) = \begin{bmatrix} \mathcal{A}(2, 1, 4, 6) \\ \mathcal{A}(2, 2, 4, 6) \\ \mathcal{A}(2, 3, 4, 6) \\ \mathcal{A}(2, 4, 4, 6) \\ \mathcal{A}(2, 5, 4, 6) \end{bmatrix}$$

is a fiber.

Slices

A slice of a tensor \mathcal{A} is a matrix obtained by fixing all but two of \mathcal{A} 's indices. For example, if $\mathcal{A} = \mathcal{A}(1:3, 1:5, 1:4, 1:7)$, then

$$\mathcal{A}(:, 3, :, 6) = \begin{bmatrix} \mathcal{A}(1, 3, 1, 6) & \mathcal{A}(1, 3, 2, 6) & \mathcal{A}(1, 3, 3, 6) & \mathcal{A}(1, 3, 4, 6) \\ \mathcal{A}(2, 3, 1, 6) & \mathcal{A}(2, 3, 2, 6) & \mathcal{A}(2, 3, 3, 6) & \mathcal{A}(2, 3, 4, 6) \\ \mathcal{A}(3, 3, 1, 6) & \mathcal{A}(3, 3, 2, 6) & \mathcal{A}(3, 3, 3, 6) & \mathcal{A}(3, 3, 4, 6) \end{bmatrix}$$

is a slice.

Where Might They Come From?

Discretization

$\mathcal{A}(i, j, k, \ell)$ might house the value of $f(w, x, y, z)$ at $(w, x, y, z) = (w_i, x_j, y_k, z_\ell)$.

Multiway Analysis

$\mathcal{A}(i, j, k, \ell)$ is a value that captures an interaction between four variables/factors.

You Have Seen them Before

Block Matrices (With Uniformly-Sized Blocks)

$$A = \begin{bmatrix} a_{11} & a_{12} & a_{13} & a_{14} & a_{15} & a_{16} \\ a_{21} & a_{22} & a_{23} & a_{24} & a_{25} & a_{26} \\ a_{31} & a_{32} & a_{33} & a_{34} & a_{35} & a_{36} \\ a_{41} & a_{42} & a_{43} & a_{44} & \boxed{a_{45}} & a_{46} \\ a_{51} & a_{52} & a_{53} & a_{54} & a_{55} & a_{56} \\ a_{61} & a_{62} & a_{63} & a_{64} & a_{65} & a_{66} \end{bmatrix}$$

Matrix entry a_{45} is the (2,1) entry of the (2,3) block:

$$a_{45} \quad \Leftrightarrow \quad \mathcal{A}(2, 3, 2, 1)$$

You Have Seen Them Before

Kronecker Products (At the Scalar Level)

$$A = \begin{bmatrix} b_{11} & b_{12} & b_{13} \\ b_{21} & b_{22} & b_{23} \\ b_{31} & b_{32} & b_{33} \end{bmatrix} \otimes \begin{bmatrix} c_{11} & c_{12} \\ c_{21} & c_{22} \end{bmatrix}$$
$$= \begin{bmatrix} b_{11}c_{11} & b_{11}c_{12} & b_{12}c_{11} & b_{12}c_{12} & b_{13}c_{11} & b_{13}c_{12} \\ b_{11}c_{21} & b_{11}c_{22} & b_{12}c_{21} & b_{12}c_{22} & b_{13}c_{21} & b_{13}c_{22} \\ b_{21}c_{11} & b_{21}c_{12} & b_{22}c_{11} & b_{22}c_{12} & b_{23}c_{11} & b_{23}c_{12} \\ b_{21}c_{21} & b_{21}c_{22} & b_{22}c_{21} & b_{22}c_{22} & b_{23}c_{21} & b_{23}c_{22} \\ b_{31}c_{11} & b_{31}c_{12} & b_{32}c_{11} & b_{32}c_{12} & b_{33}c_{11} & b_{33}c_{12} \\ b_{31}c_{21} & b_{31}c_{22} & b_{32}c_{21} & b_{32}c_{22} & b_{33}c_{21} & b_{33}c_{22} \end{bmatrix}$$

Matrix A is an **unfolding** of tensor \mathcal{A} where $\mathcal{A}(p, q, r, s) = b_{pq}c_{rs}$.

You Have Seen Them Before

Kronecker Products (At the Block Level)

$$A = \begin{bmatrix} b_{11} & b_{12} & b_{13} \\ b_{21} & b_{22} & b_{23} \\ b_{31} & b_{32} & b_{33} \end{bmatrix} \otimes \begin{bmatrix} c_{11} & c_{12} \\ c_{21} & c_{22} \end{bmatrix}$$
$$= \begin{bmatrix} b_{11}C & b_{12}C & b_{13}C \\ b_{21}C & b_{22}C & b_{23}C \\ b_{31}C & b_{32}C & b_{33}C \end{bmatrix}$$

Matrix A is a block matrix whose ij block is $b_{ij}C$.

You Have Seen Them Before

Matrix: $A = B \otimes C \otimes D$

$$A = \begin{bmatrix} b_{11} & b_{12} \\ b_{21} & b_{22} \end{bmatrix} \otimes \begin{bmatrix} c_{11} & c_{12} & c_{13} & c_{14} \\ c_{21} & c_{22} & c_{23} & c_{24} \\ c_{31} & c_{32} & c_{33} & c_{34} \\ c_{41} & c_{42} & c_{43} & c_{44} \end{bmatrix} \otimes \begin{bmatrix} d_{11} & d_{12} & d_{13} \\ d_{21} & d_{22} & d_{23} \\ d_{31} & d_{32} & d_{33} \end{bmatrix}$$

Hierarchy: A is a 2-by-2 block matrix whose entries are 4-by-4 block matrices whose entries are 3-by-3 matrices.

Tensor: $\mathcal{A} = \mathcal{D} \circ \mathcal{C} \circ \mathcal{B}$

$$\mathcal{A}(i_1, i_2, i_3, i_4, i_5, i_6) = \mathcal{D}(i_1, i_2)\mathcal{C}(i_3, i_4)\mathcal{B}(i_5, i_6)$$

**Let's look at the connection between
Kronecker products and tensors
when symmetry is present.**

$A = B \otimes C$ with Symmetric B and C

$$A = B \otimes C = \begin{bmatrix} 1 & 2 & 3 \\ 2 & 4 & 5 \\ 3 & 5 & 6 \end{bmatrix} \otimes \begin{bmatrix} 11 & 12 & 13 \\ 12 & 14 & 15 \\ 13 & 15 & 16 \end{bmatrix}$$

$$= \begin{bmatrix} 11 & 12 & 13 & | & 22 & 24 & 26 & | & 33 & 36 & 39 \\ 12 & 14 & 15 & | & 24 & 28 & 30 & | & 36 & 42 & 45 \\ 13 & 15 & 16 & | & 26 & 30 & 32 & | & 39 & 45 & 48 \\ \hline 22 & 24 & 26 & | & 44 & 48 & 52 & | & 55 & 60 & 65 \\ 24 & 28 & 30 & | & 48 & 56 & 60 & | & 60 & 70 & 75 \\ 26 & 30 & 32 & | & 52 & 60 & 64 & | & 65 & 75 & 80 \\ \hline 33 & 36 & 39 & | & 55 & 60 & 65 & | & 66 & 72 & 78 \\ 36 & 42 & 45 & | & 60 & 70 & 75 & | & 72 & 84 & 90 \\ 39 & 45 & 48 & | & 65 & 75 & 80 & | & 78 & 90 & 96 \end{bmatrix}$$

$A = B \otimes C$ with Symmetric B and C

$$A = B \otimes C = \begin{bmatrix} 1 & 2 & 3 \\ 2 & 4 & 5 \\ 3 & 5 & 6 \end{bmatrix} \otimes \begin{bmatrix} 11 & 12 & 13 \\ 12 & 14 & 15 \\ 13 & 15 & 16 \end{bmatrix}$$
$$= \begin{bmatrix} 11 & 12 & 13 & 22 & 24 & 26 & 33 & 36 & 39 \\ 12 & 14 & 15 & 24 & 28 & 30 & 36 & 42 & 45 \\ 13 & 15 & 16 & 26 & 30 & 32 & 39 & 45 & 48 \\ \hline 22 & 24 & 26 & 44 & 48 & 52 & 55 & 60 & 65 \\ 24 & 28 & 30 & 48 & 56 & 60 & 60 & 70 & 75 \\ 26 & 30 & 32 & 52 & 60 & 64 & 65 & 75 & 80 \\ \hline 33 & 36 & 39 & 55 & 60 & 65 & 66 & 72 & 78 \\ 36 & 42 & 45 & 60 & 70 & 75 & 72 & 84 & 90 \\ 39 & 45 & 48 & 65 & 75 & 80 & 78 & 90 & 96 \end{bmatrix}$$

Each block is symmetric.

$A = B \otimes C$ with Symmetric B and C

$$A = B \otimes C = \begin{bmatrix} 1 & 2 & 3 \\ 2 & 4 & 5 \\ 3 & 5 & 6 \end{bmatrix} \otimes \begin{bmatrix} 11 & 12 & 13 \\ 12 & 14 & 15 \\ 13 & 15 & 16 \end{bmatrix}$$

$$= \begin{bmatrix} 11 & 12 & 13 & 22 & 24 & 26 & 33 & 36 & 39 \\ 12 & 14 & 15 & 24 & 28 & 30 & 36 & 42 & 45 \\ 13 & 15 & 16 & 26 & 30 & 32 & 39 & 45 & 48 \\ \hline 22 & 24 & 26 & 44 & 48 & 52 & 55 & 60 & 65 \\ 24 & 28 & 30 & 48 & 56 & 60 & 60 & 70 & 75 \\ 26 & 30 & 32 & 52 & 60 & 64 & 65 & 75 & 80 \\ \hline 33 & 36 & 39 & 55 & 60 & 65 & 66 & 72 & 78 \\ 36 & 42 & 45 & 60 & 70 & 75 & 72 & 84 & 90 \\ 39 & 45 & 48 & 65 & 75 & 80 & 78 & 90 & 96 \end{bmatrix}$$

Block (i,j) equals Block (j,i)

$A = B \otimes C$ with Symmetric B and C

$$A = B \otimes C = \begin{bmatrix} 1 & 2 & 3 \\ 2 & 4 & 5 \\ 3 & 5 & 6 \end{bmatrix} \otimes \begin{bmatrix} 11 & 12 & 13 \\ 12 & 14 & 15 \\ 13 & 15 & 16 \end{bmatrix}$$
$$= \begin{bmatrix} 11 & 12 & 13 & 22 & 24 & 26 & 33 & 36 & 39 \\ 12 & 14 & 15 & 24 & 28 & 30 & 36 & 42 & 45 \\ 13 & 15 & 16 & 26 & 30 & 32 & 39 & 45 & 48 \\ \hline 22 & 24 & 26 & 44 & 48 & 52 & 55 & 60 & 65 \\ 24 & 28 & 30 & 48 & 56 & 60 & 60 & 70 & 75 \\ 26 & 30 & 32 & 52 & 60 & 64 & 65 & 75 & 80 \\ \hline 33 & 36 & 39 & 55 & 60 & 65 & 66 & 72 & 78 \\ 36 & 42 & 45 & 60 & 70 & 75 & 72 & 84 & 90 \\ 39 & 45 & 48 & 65 & 75 & 80 & 78 & 90 & 96 \end{bmatrix}$$

$$\text{If } \mathcal{A}(p, q, r, s) = b_{pq}c_{rs} \text{ then } \mathcal{A}(p, q, r, s) = \begin{cases} \mathcal{A}(q, p, r, s) \\ \mathcal{A}(p, q, s, r) \end{cases}$$

$A = B \otimes B$ with Symmetric B

$$A = B \otimes B = \begin{bmatrix} 4 & 5 & 6 \\ 5 & 7 & 8 \\ 6 & 8 & 9 \end{bmatrix} \otimes \begin{bmatrix} 4 & 5 & 6 \\ 5 & 7 & 8 \\ 6 & 8 & 9 \end{bmatrix}$$

$$= \begin{bmatrix} 16 & 20 & 24 & | & 20 & 25 & 30 & | & 24 & 30 & 36 \\ 20 & 28 & \boxed{32} & | & 25 & 35 & \boxed{40} & | & 30 & 42 & \boxed{48} \\ 24 & 32 & 36 & | & 30 & 40 & 45 & | & 36 & 48 & 54 \\ \hline 20 & 25 & 30 & | & 28 & 35 & 42 & | & 32 & 40 & 48 \\ 25 & 35 & \boxed{40} & | & 35 & 49 & \boxed{56} & | & 40 & 56 & \boxed{64} \\ 30 & 40 & 45 & | & 42 & 56 & 63 & | & 48 & 64 & 72 \\ \hline 24 & 30 & 36 & | & 32 & 40 & 48 & | & 36 & 45 & 54 \\ 30 & 42 & \boxed{48} & | & 40 & 56 & \boxed{64} & | & 45 & 63 & \boxed{72} \\ 36 & 48 & 54 & | & 48 & 64 & 72 & | & 54 & 72 & 81 \end{bmatrix}$$

$$\text{Block}(i,j) = A(i:n:n^2, j:n:n^2)$$

$$\text{Block}(2,3) = A(2:3:9, 3:3:9)$$

$A = B \otimes B$ with Symmetric B

$$A = B \otimes B = \begin{bmatrix} 4 & 5 & 6 \\ 5 & 7 & 8 \\ 6 & 8 & 9 \end{bmatrix} \otimes \begin{bmatrix} 4 & 5 & 6 \\ 5 & 7 & 8 \\ 6 & 8 & 9 \end{bmatrix}$$

$$= \begin{bmatrix} 16 & 20 & 24 & 20 & 25 & 30 & 24 & 30 & 36 \\ 20 & 28 & 32 & 25 & 35 & 40 & 30 & 42 & 48 \\ 24 & 32 & 36 & 30 & 40 & 45 & 36 & 48 & 54 \\ \hline 20 & 25 & 30 & 28 & 35 & 42 & 32 & 40 & 48 \\ 25 & 35 & 40 & 35 & 49 & 56 & 40 & 56 & 64 \\ 30 & 40 & 45 & 42 & 56 & 63 & 48 & 64 & 72 \\ \hline 24 & 30 & 36 & 32 & 40 & 48 & 36 & 45 & 54 \\ 30 & 42 & 48 & 40 & 56 & 64 & 45 & 63 & 72 \\ 36 & 48 & 54 & 48 & 64 & 72 & 54 & 72 & 81 \end{bmatrix}$$

$$\text{If } \mathcal{A}(p, q, r, s) = b_{pq}b_{rs} \text{ then } \mathcal{A}(p, q, r, s) = \begin{cases} \mathcal{A}(q, p, r, s) \\ \mathcal{A}(p, q, s, r) \\ \mathcal{A}(r, s, p, q) \end{cases}$$

A First Look at Tensor Symmetry

For a matrix, there is only one type of symmetry:

$$A(p, q) = A(q, p)$$

For an order- d tensor, there are $d! - 1$ possibilities:

$$\mathcal{A}(p, q, r, s) = \begin{cases} \mathcal{A}(q, p, r, s) \\ \mathcal{A}(r, q, p, r) \\ \mathcal{A}(s, q, r, p) \\ \vdots \end{cases}$$

Next, let's look at the connection between Kronecker products and tensors in the rank-1 setting.

Rank-1 Reshaping

If u and v are vectors, then $A = uv^T$ is a Rank-1 Matrix

$$A = \begin{bmatrix} u_1 \\ u_2 \\ u_3 \end{bmatrix} \begin{bmatrix} v_1 \\ v_2 \end{bmatrix}^T = \begin{bmatrix} u_1 v_1 & u_1 v_2 \\ u_2 v_1 & u_2 v_2 \\ u_3 v_1 & u_3 v_2 \end{bmatrix}$$

A is a rank-1 matrix

$$A = uv^T \Rightarrow \text{vec}(A) = v \otimes u$$

$$\begin{bmatrix} a_{11} \\ a_{21} \\ a_{31} \\ a_{12} \\ a_{22} \\ a_{32} \end{bmatrix} = \begin{bmatrix} u_1 v_1 \\ u_2 v_1 \\ u_3 v_1 \\ u_1 v_2 \\ u_2 v_2 \\ u_3 v_2 \end{bmatrix} = \begin{bmatrix} v_1 \\ v_2 \end{bmatrix} \otimes \begin{bmatrix} u_1 \\ u_2 \\ u_3 \end{bmatrix}$$

In the Language of Tensor Products

If u and v are vectors then $\mathcal{A} = u \circ v$ is a Rank-1 Tensor

$$\mathcal{A}(i_1, i_2) = u(i_1)v(i_2)$$

$$\mathcal{A} = u \circ v = \begin{bmatrix} u_1 \\ u_2 \\ u_3 \end{bmatrix} \circ \begin{bmatrix} v_1 \\ v_2 \end{bmatrix} \Leftrightarrow \text{vec}(\mathcal{A}) = \begin{bmatrix} u_1 v_1 \\ u_2 v_1 \\ u_3 v_1 \\ u_1 v_2 \\ u_2 v_2 \\ u_3 v_2 \end{bmatrix}$$

Higher-Order Rank-1 Tensors

If u , v , and w are vectors, then $\mathcal{A} = u \circ v \circ w$ is a Rank-1 Tensor

$$\mathcal{A}(p, q, r) = u_p v_q w_r$$

$$\mathcal{A} = u \circ v \circ w = \begin{bmatrix} u_1 \\ u_2 \end{bmatrix} \circ \begin{bmatrix} v_1 \\ v_2 \end{bmatrix} \circ \begin{bmatrix} w_1 \\ w_2 \end{bmatrix} \Rightarrow \text{vec}(\mathcal{A}) = \begin{bmatrix} u_1 v_1 w_1 \\ u_2 v_1 w_1 \\ u_1 v_2 w_1 \\ u_2 v_2 w_1 \\ u_1 v_1 w_2 \\ u_2 v_1 w_2 \\ u_1 v_2 w_2 \\ u_2 v_2 w_2 \end{bmatrix}$$

A tensor product of d vectors produces an order- d rank-1 tensor.

A Notation Detail: u-v-w versus w-v-u

$$\text{vec}(u \circ v \circ w) \equiv \begin{bmatrix} u_1 v_1 w_1 \\ u_2 v_1 w_1 \\ u_1 v_2 w_1 \\ u_2 v_2 w_1 \\ u_1 v_1 w_2 \\ u_2 v_1 w_2 \\ u_1 v_2 w_2 \\ u_2 v_2 w_2 \end{bmatrix} = w \otimes v \otimes u$$

**Let's look at how we might compute the
the nearest rank-1 tensor to
a given tensor.**

The Nearest Rank-1 Problem for Matrices

Formulation:

Given $A \in \mathbb{R}^{m \times n}$, find unit-2 norm vectors $u \in \mathbb{R}^m$ and $v \in \mathbb{R}^n$ and a nonnegative scalar σ that minimizes

$$\phi(\sigma, u, v) = \|A - \sigma uv^T\|_F.$$

SVD Solution:

If $U^T A V = \Sigma = \text{diag}(\sigma_i)$ where

$$U = [u_1 \mid \cdots \mid u_m] \quad V = [v_1 \mid \cdots \mid v_n]$$

are orthogonal and $\sigma_1 \geq \cdots \geq \sigma_n \geq 0$, then $\sigma_{\text{opt}} u_{\text{opt}} v_{\text{opt}}^T = \sigma_1 u_1 v_1^T$.

The Nearest Rank-1 Problem for Matrices

An Alternating Least Squares Approach

v = unit vector

Repeat Until Happy:

% Fix v and choose σ and u to minimize $\|A - \sigma uv^T\|_F$

$x = Av$; $\sigma = \|x\|$; $u = x/\sigma$

% Fix u and choose σ and v to minimize $\|A - \sigma uv^T\|_F$

$x = A^T u$; $\sigma = \|x\|$; $v = x/\sigma$

$\sigma_{\text{opt}} = \sigma$; $u_{\text{opt}} = u$; $v_{\text{opt}} = v$

$$\|A - \sigma uv^T\|_F^2 = \text{trace}(A^T A) - 2\sigma u^T Av + \sigma^2$$

The best u is in the direction of Av . The best v is in the direction of $A^T u$.

The Nearest Rank-1 Problem for Matrices

An Alternating Least Squares Approach

v = unit vector

Repeat Until Happy:

% Fix v and choose σ and u to minimize $\|A - \sigma uv^T\|_F$

$x = Av$; $\sigma = \|x\|$; $u = x/\sigma$

% Fix u and choose σ and v to minimize $\|A - \sigma uv^T\|_F$

$x = A^T u$; $\sigma = \|x\|$; $v = x/\sigma$

$\sigma_{\text{opt}} = \sigma$; $u_{\text{opt}} = u$; $v_{\text{opt}} = v$

This is just the power method applied to $A^T A$:

$$x = (A^T A)v, \quad v = x/\|x\|$$

Formulation

Given $\mathcal{A} \in \mathbb{R}^{m \times n \times p}$, determine unit vectors $u \in \mathbb{R}^m$, $v \in \mathbb{R}^n$, and $w \in \mathbb{R}^p$ and scalar σ so that the following is minimized:

$$\begin{aligned} \|\mathcal{A} - \sigma \cdot w \circ v \circ u\|_F &= \left(\sum_{i=1}^m \sum_{j=1}^n \sum_{k=1}^p (a_{ijk} - u_i v_j w_k)^2 \right)^{1/2} \\ &= \|\text{vec}(\mathcal{A}) - \sigma \cdot w \otimes v \otimes u\|_2 \end{aligned}$$

Nearest Rank-1 Problem for Tensors

Alternating Least Squares Framework for \min
 $\| \text{vec}(\mathcal{A}) - \sigma \cdot w \otimes v \otimes u \|_2$

v and w given unit vectors

Repeat Until Happy

Determine $x \in \mathbb{R}^m$ that minimizes $\| \text{vec}(\mathcal{A}) - w \otimes v \otimes x \|_2$

and set $\sigma = \| x \|$ and $u = x/\sigma$

Determine $y \in \mathbb{R}^n$ that minimizes $\| \text{vec}(\mathcal{A}) - w \otimes y \otimes u \|_2$

and set $\sigma = \| y \|$ and $v = y/\sigma$

Determine $z \in \mathbb{R}^p$ that minimizes $\| \text{vec}(\mathcal{A}) - z \otimes v \otimes u \|_2$

and set $\sigma = \| z \|$ and $w = z/\sigma$

Details in next Lecture. For now, we look at the special structure of these linear least square problems for the case $m = n = p = 2$.

The Nearest Rank-1 Problem for Tensors

The Case $m = n = p = 2$

$$\text{minimize} \left\| \begin{bmatrix} a_{111} \\ a_{211} \\ a_{121} \\ a_{221} \\ a_{112} \\ a_{212} \\ a_{122} \\ a_{222} \end{bmatrix} - \sigma \cdot w \otimes v \otimes u \right\|_2$$

$$u = \begin{bmatrix} \cos(\theta_1) \\ \sin(\theta_1) \end{bmatrix} = \begin{bmatrix} c_1 \\ s_1 \end{bmatrix} \quad v = \begin{bmatrix} \cos(\theta_2) \\ \sin(\theta_2) \end{bmatrix} = \begin{bmatrix} c_2 \\ s_2 \end{bmatrix} \quad w = \begin{bmatrix} \cos(\theta_3) \\ \sin(\theta_3) \end{bmatrix} = \begin{bmatrix} c_3 \\ s_3 \end{bmatrix}$$

A Highly Structured Nonlinear Optimization Problem

It Depends on Four Parameters...

$$\begin{aligned} \phi(\sigma, \theta_1, \theta_2, \theta_3) &= \left\| a - \sigma \begin{bmatrix} \cos(\theta_3) \\ \sin(\theta_3) \end{bmatrix} \otimes \begin{bmatrix} \cos(\theta_2) \\ \sin(\theta_2) \end{bmatrix} \otimes \begin{bmatrix} \cos(\theta_1) \\ \sin(\theta_1) \end{bmatrix} \right\|_2 \\ &= \left\| \begin{bmatrix} a_{111} \\ a_{211} \\ a_{121} \\ a_{221} \\ a_{112} \\ a_{212} \\ a_{122} \\ a_{222} \end{bmatrix} - \sigma \cdot \begin{bmatrix} c_3 c_2 c_1 \\ c_3 c_2 s_1 \\ c_3 s_2 c_1 \\ c_3 s_2 s_1 \\ s_3 c_2 c_1 \\ s_3 c_2 s_1 \\ s_3 s_2 c_1 \\ s_3 s_2 s_1 \end{bmatrix} \right\|_2 \end{aligned}$$

A Highly Structured Nonlinear Optimization Problem

Set $x_1 = \sigma \cos(\theta_1)$ and $y_1 = \sigma \sin(\theta_1)$ and then Reshape...

$$\phi = \left\| \begin{bmatrix} a_{111} \\ a_{211} \\ a_{121} \\ a_{221} \\ a_{112} \\ a_{212} \\ a_{122} \\ a_{222} \end{bmatrix} - \sigma \cdot \begin{bmatrix} c_3 c_2 c_1 \\ c_3 c_2 s_1 \\ c_3 s_2 c_1 \\ c_3 s_2 s_1 \\ s_3 c_2 c_1 \\ s_3 c_2 s_1 \\ s_3 s_2 c_1 \\ s_3 s_2 s_1 \end{bmatrix} \right\|_2 = \left\| \begin{bmatrix} a_{111} \\ a_{211} \\ a_{121} \\ a_{221} \\ a_{112} \\ a_{212} \\ a_{122} \\ a_{222} \end{bmatrix} - \begin{bmatrix} c_3 c_2 & 0 \\ 0 & c_3 c_2 \\ c_3 s_2 & 0 \\ 0 & c_3 s_2 \\ s_3 c_2 & 0 \\ 0 & s_3 c_2 \\ s_3 s_2 & 0 \\ 0 & s_3 s_2 \end{bmatrix} \begin{bmatrix} x_1 \\ y_1 \end{bmatrix} \right\|_2$$

This is an ordinary linear least squares problem for x_1 and y_1 if we "freeze" θ_2 and θ_3 . Solve and update σ and u_1 using

$$\begin{bmatrix} x_1 \\ y_1 \end{bmatrix} = \sigma u_1 \quad \sigma = \sqrt{x_1^2 + y_1^2}$$

A Highly Structured Nonlinear Optimization Problem

Set $x_2 = \sigma \cos(\theta_2)$ and $y_2 = \sigma \sin(\theta_2)$ and then Reshape...

$$\phi = \left\| \begin{bmatrix} a_{111} \\ a_{211} \\ a_{121} \\ a_{221} \\ a_{112} \\ a_{212} \\ a_{122} \\ a_{222} \end{bmatrix} - \sigma \cdot \begin{bmatrix} c_3 c_2 c_1 \\ c_3 c_2 s_1 \\ c_3 s_2 c_1 \\ c_3 s_2 s_1 \\ s_3 c_2 c_1 \\ s_3 c_2 s_1 \\ s_3 s_2 c_1 \\ s_3 s_2 s_1 \end{bmatrix} \right\|_2 = \left\| \begin{bmatrix} a_{111} \\ a_{211} \\ a_{121} \\ a_{221} \\ a_{112} \\ a_{212} \\ a_{122} \\ a_{222} \end{bmatrix} - \begin{bmatrix} c_3 c_1 & 0 \\ c_3 s_1 & 0 \\ 0 & c_3 c_1 \\ 0 & c_3 s_1 \\ s_3 c_1 & 0 \\ s_3 s_1 & 0 \\ 0 & s_3 c_1 \\ 0 & s_3 s_1 \end{bmatrix} \begin{bmatrix} x_2 \\ y_2 \end{bmatrix} \right\|_2$$

This is an ordinary linear least squares problem for x_2 and y_2 if we "freeze" θ_1 and θ_3 . Solve and update σ and u_2 using

$$\begin{bmatrix} x_2 \\ y_2 \end{bmatrix} = \sigma u_2 \quad \sigma = \sqrt{x_2^2 + y_2^2}$$

A Highly Structured Nonlinear Optimization Problem

Set $x_3 = \sigma \cos(\theta_3)$ and $y_3 = \sigma \sin(\theta_3)$ and then Reshape...

$$\phi = \left\| \begin{bmatrix} a_{111} \\ a_{211} \\ a_{121} \\ a_{221} \\ a_{112} \\ a_{212} \\ a_{122} \\ a_{222} \end{bmatrix} - \sigma \cdot \begin{bmatrix} c_3 c_2 c_1 \\ c_3 c_2 s_1 \\ c_3 s_2 c_1 \\ c_3 s_2 s_1 \\ s_3 c_2 c_1 \\ s_3 c_2 s_1 \\ s_3 s_2 c_1 \\ s_3 s_2 s_1 \end{bmatrix} \right\|_2 = \left\| \begin{bmatrix} a_{111} \\ a_{211} \\ a_{121} \\ a_{221} \\ a_{112} \\ a_{212} \\ a_{122} \\ a_{222} \end{bmatrix} - \begin{bmatrix} c_2 c_1 & 0 \\ c_2 s_1 & 0 \\ s_2 c_1 & 0 \\ s_2 s_1 & 0 \\ 0 & c_2 s_1 \\ 0 & c_2 s_1 \\ 0 & s_2 c_1 \\ 0 & s_2 s_1 \end{bmatrix} \begin{bmatrix} x_3 \\ y_3 \end{bmatrix} \right\|_2$$

This is an ordinary linear least squares problem for x_3 and y_3 if we "freeze" θ_1 and θ_2 . Solve and update σ and u_3 using

$$\begin{bmatrix} x_3 \\ y_3 \end{bmatrix} = \sigma u_3 \quad \sigma = \sqrt{x_3^2 + y_3^2}$$

A Common Framework for Tensor-Related Optimization

- Choose a subset of the unknowns such that if they are (temporarily) fixed, then we are presented with some standard matrix problem in the remaining unknowns.
- By choosing different subsets, cycle through all the unknowns.
- Repeat until converged.

*In tensor computations, the “standard matrix problem” that we end up solving is usually the linear least squares problem. In that case, the overall solution process is referred to as **alternating least squares**.*

Optional “Fun” Problems

Problem E1. Consider the the three linear least (LS) squares problems that arise when the alternating least squares framework is applied to the 2-by-2-by-2 problem. Outline a solution approach when these linear LS problems are solved using the method of normal equations. (Recall that the method of normal equations for the LS problem $\min \|Mu - b\|_2$ involves solving the symmetric positive definite linear system $M^T M u = M^T b$.)

Problem A1. Repeat E1 but when $\mathcal{A} \in \mathbb{R}^{2 \times 2 \times \dots \times 2}$ is an order-d tensor.

Closing Remarks

Where Do We Go From Here?

To sums of rank-1's...

$$\text{vec}(\mathcal{A}) = \sum_{k=1}^r \sigma_k w_k \otimes v_k \otimes u_k$$

To more general unfoldings...

$$\mathcal{A} \in \mathbb{R}^{4 \times 2 \times 3} \Rightarrow \begin{bmatrix} a_{111} & a_{121} & a_{131} & a_{112} & a_{122} & a_{132} \\ a_{211} & a_{221} & a_{231} & a_{212} & a_{222} & a_{232} \\ a_{311} & a_{321} & a_{331} & a_{312} & a_{322} & a_{332} \\ a_{411} & a_{421} & a_{431} & a_{412} & a_{422} & a_{432} \end{bmatrix}$$

To more complicated multilinear optimizations...

$$\begin{aligned} & \min && \|\text{vec}(\mathcal{A}) - (W \otimes V \otimes U)s\|_2 \\ & U, V, W \in \mathbb{R}^{n \times n} \text{ orthogonal} \\ & s \in \mathbb{R}^3 \end{aligned}$$

How Will the Structured Matrix Computations Show Up?

Tensor computations are typically disguised matrix computations and that is because of

Kronecker Products

$$A = A_1 \otimes A_2 \otimes A_3 \quad \text{an order 6 tensor}$$

Tensor Unfoldings

$$\text{Rubik Cube} \quad \longrightarrow \quad 3 \times 9 \quad \text{matrix}$$

Alternating Least Squares

Multilinear optimization via component-wise linear optimization

These are the three ways that structured tensor computations will lead to structured matrix computations.

Preparation for the Next Big Thing...

Scalar-Level Thinking

1960's ↓



The factorization paradigm:
 LU , LDL^T , QR , $U\Sigma V^T$, etc.

Matrix-Level Thinking

1980's ↓



Cache utilization, parallel
computing, LAPACK, etc.

Block Matrix-Level Thinking

2000's ↓



New applications, factoriza-
tions, data structures, non-
linear analysis, optimization
strategies, etc.

Tensor-Level Thinking

A Changing Definition of “Big”

In Matrix Computations, to say that $A \in \mathbb{R}^{n_1 \times n_2}$ is “big” is to say that both n_1 and n_2 are big.

In Tensor Computations, to say that $\mathcal{A} \in \mathbb{R}^{n_1 \times \dots \times n_d}$ is “big” is to say that $n_1 n_2 \dots n_d$ is big and this need not require big n_k . E.g. $n_1 = n_2 = \dots = n_{1000} = 2$.

Algorithms that scale with d will induce a transition...

Matrix-Based Scientific Computation



Tensor-Based Scientific Computation